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User's Manual

RA78K0S Ver. 2.00

Assembler Package

Structured Assembler Language

Target Device

78K0S Microcontrollers

Document No. U17389EJ2V0UM00 (2nd edition)

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INTRODUCTION

This manual has been written to help users obtain an accurate understanding of the coding method used for the structured assembler preprocessor (hereafter referred to as the "structured assembler") that is included in the RA78K0S Microcontrollers Assembler Package (hereafter referred to as "the RA78K0S").

This manual does not explain methods for using programs other than the structured assembler nor does it describe structured assembler operation methods.

Therefore, before reading this manual, read the **RA78K0S Assembler Package User's Manual LANGUAGE (U17390E)** and **OPERATION (U17391E)**.

The contents of this manual are intended for use with Ver. 1.50 of the RA78K0S.

[Target Readers]

The RA78K0S is intended for users who understand the functions and instructions of the microcontroller to be developed (78K0S Microcontrollers).

Readers requiring a description of the 78K0S Microcontrollers should refer to the target chip's User's Manual.

[Organization]

This manual consists of the following chapters and appendixes:

CHAPTER 1 GENERAL

This chapter describes the functions (the role, etc.) of the structured assembler in software development for microcontrollers.

CHAPTER 2 SOURCE PROGRAM CODING METHODS

This chapter describes methods for source program configuration, coding syntax, and other principal rules and conventions concerning the coding of source programs.

CHAPTER 3 CONTROL STATEMENTS

Control statements are used to describe the "if~else~endif" indicators of the program structure. This chapter describes control statement functions and coding methods.

CHAPTER 4 EXPRESSIONS

Assignments and arithmetic operations are entered as expressions. This chapter describes expression functions and coding methods.

CHAPTER 5 DIRECTIVES

This chapter presents use examples in describing how to write and use structured assembler directives.

CHAPTER 6 CONTROL INSTRUCTIONS

This chapter presents use examples in describing how to write and use structured assembler control instructions.

APPENDIX A SYNTAX LISTS

This appendix presents a structured assembler syntax list.

APPENDIX B LISTS OF GENERATED INSTRUCTIONS

This appendix presents a list of instructions generated by the structured assembler.

The instruction sets are not detailed in this manual.

For these instructions, refer to the user's manual of the microcontroller to be developed.

[How to Read This Manual]

Those using an assembler for the first time are encouraged to read from **CHAPTER 1 GENERAL** of this manual. Those who have a general understanding of assembler programs may skip this chapter. However, all readers should read section “**1.3 Before Starting Program Development**”.

[Conventions]

The following symbols and abbreviations are used throughout this manual:

- : : Indicates that the same expression is repeated.
- []: Item(s) in brackets can be omitted.
- ' ': Characters enclosed in ' ' (quotation marks) will be listed as they appear.
- " ": Characters enclosed in " " (double quotation marks) are titles of chapters, paragraphs, sections, diagrams or tables to which the reader is asked to refer.
- ___: Indicates an important point, or characters that are to be input in a usage example.
- : Indicates one blank space.
- Δ: Indicates one or more blank or TAB.
- ∇: Indicates zero or more blanks or TABs (i.e. blanks may be omitted).
- /: Indicates a break between characters.
- ~: Indicates continuity.
- [↵]: Indicates pressing of the Return key.
- Note:** Footnote for item marked with **Note** in the text
- Caution:** Information requiring particular attention
- Remark:** Supplementary information

[Related Documents]

The documents related to this manual are listed below.

The related documents indicated in this publication may include preliminary versions. However, preliminary versions are not marked as such.

Document related to development tools (user's manuals)

Document Name		Document Number
CC78K0S Ver. 2.00 C Compiler	Operation	U17416E
	Language	U17415E
RA78K0S Ver. 2.00 Assembler Package	Operation	U17391E
	Language	U17390E
	Structured Assembly Language	This manual
SM+ System Simulator	Operation	U18601E
	User Open Interface	U18602E
SM78K Series Ver.2.52 System Simulator	Operation	U16768E
ID78K0S-NS Integrated Debugger Ver. 2.52	Operation	U16584E
ID78K0S-QB Integrated Debugger Ver. 3.00	Operation	U17287E
PM+ Ver. 6.30 Project Manager		U18416E

Caution The contents of the above related documents are subject to change without notice. Be sure to use the latest edition of each document when designing your system.

[MEMO]

CONTENTS

CHAPTER 1	GENERAL ...	13
1.1	Overview ...	13
1.2	Overview of Functions ...	14
1.2.1	Main Functions ...	14
1.2.2	Flowchart of Program Development ...	15
1.3	Before Starting Program Development ...	16
1.3.1	Maximum performance ...	16
1.3.2	Word symbols and byte symbols ...	17
1.3.3	Definition of label ...	17
CHAPTER 2	SOURCE PROGRAM CODING METHODS ...	18
2.1	Basic Configuration ...	18
2.2	Elements ...	20
2.3	Reserved Words ...	23
2.4	Label Generation Rules ...	25
2.5	Size Specification ...	26
2.6	Data Sizes ...	27
2.7	Comments ...	29
2.8	Tool Information ...	30
2.9	Output Results of Input Source Files by ST78K0S ...	31
CHAPTER 3	CONTROL STATEMENTS ...	32
3.1	Control Statement Characters ...	32
3.2	Nesting ...	33
3.3	Register Specification ...	34
3.4	Control Statement Functions ...	36
3.4.1	Conditional branch ...	37
	Conditional branch if ...	37
	Conditional branch if_bit ...	40
	Conditional branch switch ...	43
3.4.2	Conditional loop ...	47
	Conditional loop for ...	47
	Conditional loop while ...	50
	Conditional loop while_bit ...	52
	Conditional loop until ...	54
	Conditional loop until_bit ...	56
	Conditional loop break ...	57
	Conditional loop continue ...	58
	Conditional loop goto ...	59
3.5	Conditional Expressions ...	60
3.5.1	Comparison expressions ...	61
	Comparison expressions Equal (==) ...	63
	Comparison expressions NotEqual (!=) ...	66
	Comparison expressions LessThan (<) ...	69
	Comparison expressions GreaterThan (>) ...	72
	Comparison expressions GreaterEqual (>=) ...	75
	Comparison expressions LessEqual (<=) ...	78
	Comparison expressions FOREVER (forever) ...	81
3.5.2	Test bit expressions ...	83
	Test bit expressions Positive logic (bit) ...	84
	Test bit expressions Negative logic (bit) ...	87
3.5.3	Logical operations ...	90
	Logical operations Logical AND (&&) ...	91
	Logical operations Logical OR () ...	94

CHAPTER 4 EXPRESSIONS ...	96
4.1 Overview ...	96
4.2 Assignment Statements ...	99
Assignment statements Assign (=) ...	99
Assignment statements IncrementAssign (+=) ...	104
Assignment statements DecrementAssign (-=) ...	107
Assignment statements LogicalANDAssign (&=) ...	110
Assignment statements LogicalORAssign (=) ...	113
Assignment statements LogicalXORAssign (^=) ...	116
Assignment statements RightShiftAssign (>>=) ...	119
Assignment statements LeftShiftAssign (<<=) ...	121
4.3 Count Statements ...	123
Count statements Increment (++) ...	123
Count statements Decrement (--) ...	125
4.4 Exchange Statements ...	127
Exchange statements Exchange (<->) ...	127
4.5 Bit Manipulation Statements ...	130
Bit manipulation statements Set bit (=) ...	130
Bit manipulation statements Clear bit (=) ...	133
 CHAPTER 5 DIRECTIVES ...	 136
5.1 Overview ...	136
5.2 Directive Functions ...	137
#define ...	138
#ifdef / #else / #endif ...	140
#include ...	142
#defcallt ...	143
 CHAPTER 6 CONTROL INSTRUCTIONS ...	 144
6.1 Overview ...	144
6.2 Assembler Control Instructions ...	145
6.3 Control Instruction Functions ...	147
\$PROCESSOR / \$PC ...	148
\$KANJICODE ...	149
 APPENDIX A SYNTAX LISTS ...	 150
 APPENDIX B LISTS OF GENERATED INSTRUCTIONS ...	 154
 APPENDIX C INDEX ...	 165

LIST OF FIGURES

Figure No. Title , Page

1-1	ST78K0S Function ...	14
1-2	Program Development Flowchart ...	15
3-1	Nesting example ...	33

LIST OF TABLES

Table No. Title Page

1-1	Maximum Performance of ST78K0S ...	16
2-1	Structured Assembly Language Coding ...	18
2-2	Alphanumeric Characters ...	20
2-3	Special Characters ...	21
2-4	Invalid Characters ...	22
2-5	Reserved Word ...	23
2-6	Data Sizes ...	27
2-7	Output by ST78K0S ...	31
3-1	List of Control statements ...	36
3-2	Generated Instructions for switch Statements ...	45
3-3	Generated Instructions for Comparison Instructions ...	61
3-4	Comparison Expressions ...	62
3-5	Test Bit Expressions ...	83
3-6	Logical Operations ...	90
3-7	Generated Instructions (Control Statement in Lower Case Letters) for Logical AND ...	91
3-8	Generated Instructions (Control Statement in Upper Case Letters) for Logical AND ...	92
3-9	Generated Instructions for Logical OR ...	94
4-1	Assignment Statements ...	96
4-2	Count Statements ...	97
4-3	Exchange Statements ...	98
4-4	Bit Manipulation Statements ...	98
4-5	Generated Instructions for Assignments ...	103
4-6	Generated Instructions for Increment Assignments ...	106
4-7	Generated Instructions for Decrement Assignments ...	109
4-8	Generated Instructions for Logical AND Assignments ...	112
4-9	Generated Instructions for Logical OR Assignments ...	115
4-10	Generated Instructions for Logical XOR Assignments ...	118
4-11	Generated Instructions for Increment ...	124
4-12	Generated Instructions for Decrement ...	126
4-13	Generated Instructions for Exchange ...	129
4-14	Generated Instructions for Set Bit ...	132
4-15	Generated Instructions for Clear Bit ...	135
5-1	List of Directives ...	137
6-1	Control Instructions that Can Be Entered Only in Module Headers ...	145
6-2	Control Instructions that Are Recognized as the Module Body ...	146
6-3	Control Instruction List ...	147
6-4	Interpretation of Kanji Code ...	149
A-1	Control Statements ...	150
A-2	Conditional Expressions ...	151
A-3	Expressions ...	152
A-4	Directives ...	153
A-5	Control Instructions ...	153
B-1	Generated Instructions for Comparison Expressions ...	154
B-2	Generated Instructions for Test Bit Expressions ...	157
B-3	Generated Instructions for Logic Expressions ...	158
B-4	Expressions ...	161

CHAPTER 1 GENERAL

This chapter describes the functions (the role, etc.) of the structured assembler preprocessor in software development for microcontrollers.

1.1 Overview

The RA78K0S structured assembler preprocessor (ST78K0S) is a program in the "RA78K0S Assembler Package" that is used for software development of microcontrollers in the 78K0S Series.

The ST78K0S converts structured assembly statements such as "if~else~endif" and "for~next" into assembly language source . Control statements are used to enter "if~else~endif" and "for~next" descriptions.

As such, the ST78K0S offers the following three advantages.

(1) Programs are easy to write

- Each program structure can be written as is, which facilitates the development process from design to coding.
- There is no need to consider label names for branching.
- Transfer instructions that contain large amounts of code can be entered as assignment statements.

(2) Programs are easy to read.

- Program structure is easy to understand.
- Operations and transfers between memory registers can be entered in a single statement.
- Other programmer's programs are easy to read.
- Program maintenance (revision) is easy.

(3) Facilitates desktop debugging

- Coding can be done on a one-to-one correspondence with the detail design, thus facilitating desktop debugging.

1.2 Overview of Functions

The ST78K0S analyzes various control statements, expressions, and directives within a structured assembler source program that are coded according to a specific language specification and outputs an assembler source that serves as an input source file for the assembler.

Structured statements can be output as comments and converted assembler instructions and ordinary assembly language can all be output as secondary source files.

Error messages are output when errors occur.

Figure 1-1 ST78K0S Function



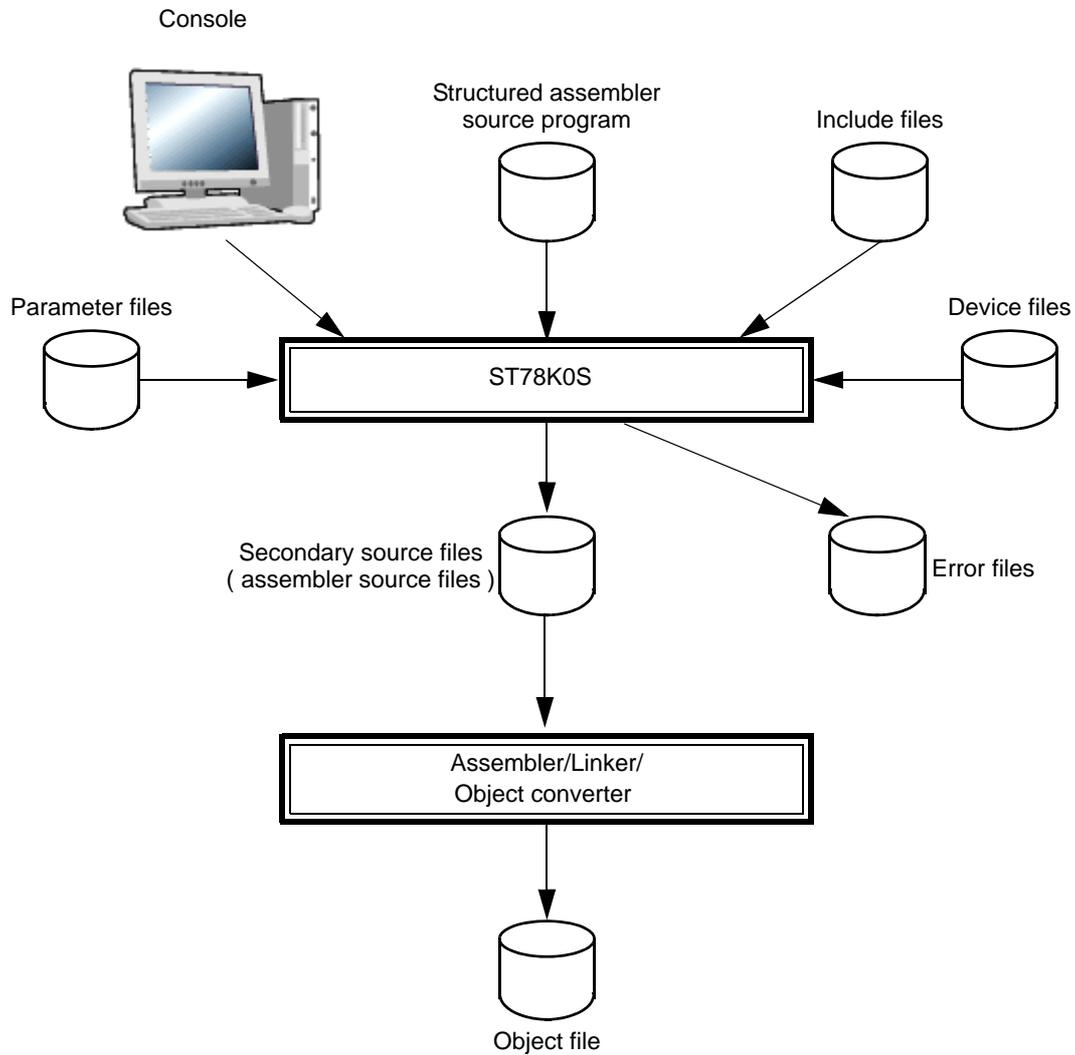
1.2.1 Main Functions

- (1) Program coding is facilitated by an abundance of C-like control statements.
- (2) C-like assignment statements and assignment operators can be used in coding.
- (3) Control structures and assignment statements can be coded for bit processing.
- (4) It includes C-like symbol definition directives, conditional processing functions, and include directives.
- (5) Since it is the preprocessor that outputs assembler source, code optimization can be performed following conversion by the ST78K0S.
- (6) Using a directive that can be converted into a CALLT instruction, routines to be registered to the CALLT table can be determined following the development of a program.
- (7) Easy-to-read assembly lists can be created by changing the assembler source output position.

1.2.2 Flowchart of Program Development

Figure 1-2 shows a flowchart of program development.

Figure 1-2 Program Development Flowchart



Remark Obtain the device file by downloading it from the Online Delivery Service (ODS), which can be accessed from the following Website.

<http://www.necel.com/micro/ods/eng/tool/DeviceFile/list.html>

1.3 Before Starting Program Development

The maximum performance of the ST78K0S and points to be noted are described below.

1.3.1 Maximum performance

Table 1-1 Maximum Performance of ST78K0S

Item	Maximum value
Line length (not including LF or CR)	2048 characters
Number of symbols registered in #define directive (excluding reserved words)	512 symbols
Character length of symbol registered in #define directive	31 characters
Nesting levels in control statement	31 levels
Nesting levels in #ifdef directive	8 levels
#defcallt directives	32
Nesting of #include directives	Not supported
Number of redefinitions by #define directive	31 times
Number of operands assigned in a series	33 (Note 1)
Logical operator operands	17 (Note 2)
Number of symbols defined by option "-d"	30
Number of include file paths specifiable by -i option	64

Notes 1. The maximum value is expressed as follows.

$$S1 = S2 = \dots S32 = S33$$

Up to 33 symbols and 32 equal (=) signs can be inserted.

Notes 2. The maximum value is expressed as follows.

$$\text{expression 1}\&\&\text{expression 2}\&\& \dots \&\&\text{expression 16}\&\&\text{expression 17}$$

Up to 17 expressions and 16 "&&" (or "||") signs can be inserted.

CHAPTER 2 SOURCE PROGRAM CODING METHODS

This chapter describes coding methods for source programs etc.

2.1 Basic Configuration

Source programs consist of structured assembly language and (pure) assembly language.

For further description of assembly language, see the RA78K0S Assembler Package Language User's Manual.

Each line (between two LFs) can contain up to 2048 characters.

The types of coding used in structured assembly language are listed below in [Table 2-1](#).

Table 2-1 Structured Assembly Language Coding

Type		Coding
ST78K0S statement	Control statement	Conditional branch if ~ elseif ~ else ~ endif if_bit ~ elseif_bit ~ else ~ endif switch ~ case ~ default ~ ends
		Conditional loop for ~ next while ~ endw while_bit ~ endw repeat ~ until repeat ~ until_bit
		Other break, continue, goto
	Expression	Assignment statement Assign (=), assignment plus operation (+=, etc.), shift (rotate) assignment (>>=, etc.)
		Count statement Increment (++), Decrement (--)
		Exchange statement Exchange (<->)
		Bit manipulation statement Set bit (=), Clear bit (=)
	Conditional expression	Comparison expression ==, !=, <, >, >=, <=
Test bit expression Bit symbol, !bit symbol		
Logical operation Logical AND (&&), Logical OR ()		

(1) Control statements

Control statements include "if ~ elseif ~ else ~ endif", "if_bit ~ elseif_bit ~ else ~ endif", and "switch ~ case ~ default ~ ends" statements that represent conditional branches, "for ~ next", "while ~ endw", "while_bit ~ endw", "repeat ~ until", and "repeat ~ until_bit" statements that represent conditional loops, and "break", "continue", and "goto" statements that represent loop exit processing. For details, see "[CHAPTER 3 CONTROL STATEMENTS](#)".

(2) Expressions

Expressions include assignment statements, count statements (increment and decrement), exchange statements, and bit manipulation statements. For details, see "[CHAPTER 4 EXPRESSIONS](#)".

(3) Conditional expressions

Conditional expressions include comparison expressions, test bit expression, and logical operations. For details, see "[3.5 Conditional Expressions](#)".

2.2 Elements

(1) Character set

Letters, numerals, and special characters can be used in source programs.

Table 2-2 Alphanumeric Characters

Name		Character
Numerals		0 1 2 3 4 5 6 7 8 9
Letters	Upper case	A B C D E F G H I J K L M N O P Q R S T U V W X Y Z
	Lower case	a b c d e f g h i j k l m n o p q r s t u v w x y z

In the ST78K0S, only the first character in control statements are case-sensitive. Any lower case letters that appear after the first character are converted to upper case letters. However, secondary source files are output using the case specifications in which they were entered.

Table 2-3 Special Characters

Character	Name	Use
?	Question mark	Character used as letter
@	Unit price symbol	Character used as letter
_	Underlining	Character used as letter
	White space	Delimiter symbol for phrases
HT	Horizontal tab	Character used as white space
,	Comma	Delimiter symbol for operands
.	Period	Bit position symbol for bit symbols
"	Double quotation mark	Specification character for #INCLUDE directive's disk-type file names
'	Single quotation mark	Symbol used to mark start and end of character constant
+	Plus symbol	Positive sign or increment operation
-	Minus symbol	Negative sign or decrement operation
&	Ampersand	Logical AND operator
	Separator symbol	Logical OR operator
^	Upward arrow symbol	Exclusive OR operator
(Left parenthesis	Change in operation sequence or expression in control statement
)	Right parenthesis	
=	Equal symbol	Assignment operator, comparison operator
:	Colon	Delimiter symbol for labels
;	Semicolon	Comment start symbol or delimiter symbol in control statement expressions
#	Sharp symbol	First character in ST78K0S directive or immediate display symbol
\$	Dollar sign	Location or counter value Display symbol in control instruction
!	Exclamation point	Direct addressing specification symbol, negation display symbol
<	Not equal (less than) symbol	Comparison operator
>	Not equal (more than) symbol	
\	Back slash	Directory specification symbol
[Left bracket	Indirect address specification symbol
]	Right bracket	
LF	Line feed	End of line symbol

An error will occur if any of the following [Invalid Characters](#) are entered.

Table 2-4 Invalid Characters

Type	ASCII code
Illegal characters	00H to 08H, 0BH, 0CH, 0EH to 1FH, 7FH
Unrecognized special characters	% (25H), ' (60H), { (7BH), } (7DH), ^ (7EH)
Other characters	80H ~ 0FFH

When an illegal character is entered, an error occurs and each illegal character is replaced by a period (.) when a secondary file is output.

However, invalid characters can be used in comments.

(2) Identifiers

Identifiers are names that are attached to numerical data, addresses, etc.

Identifiers are used to make the contents of source programs easier to identify.

Use #define statements to define details of identifiers (see also "[5.2 Directive Functions](#)").

(3) Symbols

The last character in the symbol name determines whether the ST78K0S generates a byte access instruction or a word access instruction. The default setting is P (pair), which can be changed via the -sc option.

All character strings other than reserved word symbols can be handled as user symbols. All alphanumeric characters and all other characters that can be established as English alphabet characters can be used as user symbols.

(4) Constants

Structured assembly language does not include any constants. However, assembly language constants can be output as is to secondary files (for details of assembly language constants, refer to the RA78K0S Assembler Package Language User's Manual.

(5) Expressions

Expressions are constants, special characters, and symbols that are combined using operators (for details of assembly language expressions, see the RA78K0S Assembler Package Language User's Manual.

Be sure to enclose in parentheses any symbols that are separated by white spaces within an assembly language expression.

< Examples >

- Coding method for assembler

```
MOV    A , # ( SYM AND 0FFH )
MOV    A , LABEL + 1
```

- Coding method for ST78K0S structured assembler source program

```
A = # ( SYM AND 0FFH )
A = ( LABEL + 1 )
```

2.3 Reserved Words

Table 2-5 lists reserved words in structured assembly language.

For information on instructions and sfr symbols, see the target device's User's Manual.

Table 2-5 Reserved Word

Type	Reserved word
Control statements	IF, IF_BIT, ELSEIF, ELSEIF_BIT, ELSE, ENDIF
	SWITCH, CASE, DEFAULT, ENDS
	FOR, NEXT
	WHILE, WHILE_BIT, ENDW
	REPEAT, UNTIL, UNTIL_BIT
	BREAK, CONTINUE, GOTO
Directives	DFINE
	IFDEF, ELSE, ENDIF
	INCLUDE
	DEFCALLT, ENDCALLT
Operators	++, --
	=, +=, -=, *=, /=, &=, =, ^=, <<=, >>=, <->
	==, !=, <, >=, >, <=, FOREVER
Assembler operators	MOD, NOT
	AND, OR, XOR
	EQ, NE, GT, GE, LT, LE
	SHL, SHR
	HIGH, LOW, BANKNUM
	DATAPOS, BITPOS, MASK

Table 2-5 Reserved Word

Type	Reserved word
Assembler control instructions	PROCESSOR, PC
	DEBUG, NODEBUG, DEBUGA, NODEBUGA, DG, NODG
	XREF, XR, NOXREF, NOXR
	TITLE, TI
	SYMLIST, NOSYMLIST
	FORMFEED, NOFORMFEED
	WIDTH, LENGTH
	TAB
	KANJICODE
	IC
	EJECT, EJ
	LIST, LI, NOLIST, NOLI
	GEN, NOGEN
	COND, NOCOND
	SUBTITLE, ST
SET, RESET	
_IF, _ELSEIF, IF, ELSEIF, ELSE, ENDIF	
Registers	CY, Z
	A, X, B, C, D, E, H, L
	R0, R1, R2, R3, R4, R5, R6, R7
	PSW
	AX, BC, DE, HL
	RP0, RP1, RP2, RP3
	SP
Other	DGS, DGL, TOL_INF, SJIS, EUC, NONE

2.4 Label Generation Rules

When using control statements in assembler language instructions, the ST78K0S generates labels for branch instructions.

Labels generated by the ST78K0S have the format "?Lddd".

The "ddd" represents a decimal value of 1 or more, output with suppression of zeros and left alignment. Therefore, do not enter any labels using this "?Lddd" format.

2.5 Size Specification

Size specifications can be made to change the data size of symbols entered in the left or right sides of an assignment expression or a conditional expression or case symbols in switch statements.

(1) Description format

(Δ size specification character Δ)

(2) Function

- If the size character is either "B" or "b", the data size is changed to bytes.

(3) Description

- An error will occur if the size specification character is incorrect.
- An error will occur if a size specification is entered in an assignment expression or a conditional expression which does not support size specifications.
- If a size specification is made to a register, coding can only be done using the same specification. The data size cannot be changed. If the data size is different, an error will occur.
- When specifying a user symbol, be sure to change the data size to the specified data size.
- If a size specification has been entered for a direct access specification symbol or an indirect access specification symbol or for immediate data, the size specification will be ignored and the data size will not be changed.
- Word access cannot be specified in size specifications.

2.6 Data Sizes

The ST78K0S checks the data size of symbols. This is because the symbols differ according to the instruction being generated. However, the ST78K0S allows the assembler to determine whether or not the symbol definitions and constants are entered correctly.

The data sizes checked by the ST78K0S are listed below.

Table 2-6 Data Sizes

Symbol in Generated Instruction Table	Description
a	CY
b	Bit symbols (except [HL]. β) This ST78K0S recognizes bit sfrs and symbols entered using the format " α , β " as bit symbols. Items that can be entered as " α " include byte user symbols, word user symbols, byte-specified user symbols, sfrs, A, PSW, and constants. Items that can be entered as " β " include byte user symbols, word user symbols, and constants.
c	[HL]. β Items that can be entered as " β " include byte user symbols, word user symbols, and constants.
d	Byte user symbols
e	Byte-specified user symbols, sfrs that overlaps saddr
f	A
g	Byte registers (except A, R0, and R1)
h	R0
i	R1
j	sfr
k	PSW
l	Word user symbols
m	sfrp that overlaps saddrp
n	AX
o	Word register (except AX and RP0)
p	RP0
q	sfrp
r	SP
s	Direct access specification symbols These are symbols that are specified using the format "!addr". Byte user symbols, word user symbols, constants, and \$ can be entered as "addr".
t	Indirect access specification symbols These are symbols that are specified using the format "[HL]" and "[HL + byte]". Byte user symbols, constants, and \$ can be entered as "byte".

Table 2-6 Data Sizes

Symbol in Generated Instruction Table	Description
u	Special indirect access specification symbols These are symbols that are specified using the format "[DE]".
v	Immediate data These are symbols that are specified using the format "#date". Byte user symbols, word user symbols, constants, and \$ can be entered as "date".

2.7 Comments

Any character string that appears after a semicolon (;) until the next line feed (LF) is regarded as a comment statement, which is not processed but is simply output to a secondary file. Comment statements can be entered at any position in a line of code.

However, since semicolons are used between parentheses as expression delimiters in the "for ~ next" syntax, the two semicolons that are entered between parentheses are not regarded as the start of a comment statement.

All of the characters listed under "2.2 (1) Character set" can be used in comments.

Processing of illegal characters does not occur when the illegal characters are included in a comment or comment statement.

2.8 Tool Information

The ST78K0S outputs tool information.

If an input source file contains tool information that has been output by the ST78K0S, the "\$" character at the start of the information is replaced with ";".

The output position is the end of the module header. The only types of statements that can be entered in module headers are assembler control instructions, comment statements, and line feeds.

(1) Output format

```
$TOL_INF      2FH , second parameter , third parameter , 0FFFFH
```

2FH indicates that it is tool information output by the ST78K0S preprocessor.

The second parameter indicates the version number of this preprocessor.

The version number is output either as a hexadecimal value or, if the value is not converted, as the decimal number image that was shown at startup.

< Example >

Version number 3.10 -> 310H

The third parameter is used to indicate this preprocessor's error messages.

```
0H :          Normal end
1H :          Fatal error, exited
2H :          Warning, exited
3H :          Fatal error and warning, exited
```

0FFFFH indicates language-related information. This is a fixed value for this preprocessor.

2.9 Output Results of Input Source Files by ST78K0S

Input source files are output as follows by the ST78K0S.

Table 2-7 Output by ST78K0S

Input source program file	Secondary source program file
ST78K0S control statements ST78K0S expression statements	Output as comments
ST78K0S directives	Not output
#INCLUDE	Outputs include contents
Source alias set by #IFDEF	Not output
Comments	Output as comments
Other lines	Output as is

CHAPTER 3 CONTROL STATEMENTS

This chapter presents examples in describing control statement functions.

Control statements are used to structurally code the flow of program control (see also "[3.4 Control Statement Functions](#)").

3.1 Control Statement Characters

The instruction generated by a control statement differs fundamentally depending on whether upper case or lower case letters are used in the control statement. For example, the different statement sizes between "if ~ endif" and "IF ~ ENDIF" can preclude direct branching via the conditional branch instruction generated by processing of the condition expression.

However, ensuring that the statement will always be branched correctly has the disadvantage of reducing the program's efficiency as an object.

As a solution to this problem, the user is able to set upper or lower case in order to improve the object efficiency rate. If there is no need to improve the object efficiency rate, the user can omit changing the character size as long as coding uses upper case letters.

Since control statements generate conditional branch instructions, be sure to specify whether or not the relative address is within 128 bytes.

In control statements, "if" and "elseif" are reserved words. The ST78K0S determines whether the first character in a control statement reserved word is an upper case or lower case letter.

IF, If : First letter is upper case, so coding is determined as upper case.

if, iF : First letter is lower case, so coding is determined as lower case.

If entered in upper case : branches using a combination of conditional branch instruction and BR directive.

If entered in lower case : branches directly using a conditional directive.

Paired control statements (such as "if, else, endif") can have mixed upper case and lower case letters. In other words, it is possible to enter one as "IF ~ else ~ ENDIF".

3.2 Nesting

Control statements can be nested. Generally, up to 31 nesting levels are allowed. However, control statements cannot be intersected.

Figure 3-1 Nesting example

< Example of incorrect coding >

```
while ( A < B )
    if ( A == #4 )
        break ;
endw
endif
```

Error occurs due to intersecting.

< Example of correct coding >

```
while ( A < B )
    if ( A == #4 )
        break ;
    endif
endw
```

"if" statement is correctly nested within "while" statement.

3.3 Register Specification

(1) Description format

([Δ] [=] [Δ] register name [Δ])

(2) Function

- If a register is specified immediately after a comparison expression

After the instruction to transfer the left side to the specified register, a comparison expression is generated to compare the specified register with the right side.

< Example >

Output source	Input source
<pre> CMP SYM1 , #5 BZ \$?L1 CMP SYM2 , #0 BC \$?L1 MOV A , SYM3 CMP A , #80H BNC \$?L1 ?L1 :</pre>	<pre> if (SYM1 != #5 && SYM2 >= #0&&SYM3 < #80H (A)) endif</pre>

- If a register is specified after a control statement

During the generated of each comparison expression, after the instruction for transferring the left side to the specified register is generated, a comparison expression is generated to compare the specified register with the right side.

< Example >

Output source	Input source
<pre> MOV A , R4 CMP A , #5 BZ \$?L2 MOV A , R2 CMP A , #0 BC \$?L2 MOV A , R3 CMP A , #80H BNC \$?L2 ?L2 :</pre>	<pre> if (R4 != #5 && R2 >= #0 && R3 <# 80H) (A) endif</pre>

- If both (a) and (b) are specified

The register specification that immediately follows a comparison expression takes priority. After the instruction for transferring the left side to the specified register is generated, a comparison expression is generated to compare the specified register with the right side.

As for an expression in which there is no register specification immediately after a comparison expression, after the instruction for transferring the left side to the specified register is generated according to the register specification following the control statement, a comparison expression is generated to compare the specified register with the right side.

< Example >

Output source	Input source
<pre> MOV A , DATA1 CMP A , #5 BZ \$?L3 MOV A , DATA2 CMP A , #0 BC \$?L3 MOV A , DATA3 CMP A , #80H BNC \$?L3 ?L3 :</pre>	<pre> if (DATA1!=#5 && DATA2 >=#0 (A) && DATA3<#80H) (A) endif</pre>

(3) Description

- Register specifications can be used in if statements, elseif statements, switch statements, for statements, while statements, and until statements. However, if the conditional expression is a bit expression, any register specified in the control statement is ignored.
- For a list of register names, see [Table 2-5](#).
sfr specifications can also be entered.
- The processing for an assignment statement within a for statement is the same as for comparison expressions.

3.4 Control Statement Functions

The following pages describe the functions of the various control statements.

The use examples show as comment statements the source files to which generated instructions are input.

Table 3-1 List of Control statements

Type	Description	Remark
Conditional branch	if ~ elseif ~ else ~ endif	
	if_bit ~ elseif_bit ~ else ~ endif	
	switch ~ case ~ default ~ ends	
Conditional loop	for ~ next	Repetition of increment specification
	while ~ endw	Repetition of conditional expression testing before processing
	while_bit ~ endw	Repetition of conditional expression testing before processing
	repeat ~ until	Repetition of conditional expression testing after processing
	repeat ~ until_bit	Repetition of conditional expression testing after processing
	break	Extraction of loop block
	continue	Repetition of loop block
	goto	Escape to go to exception processing

3.4.1 Conditional branch

Conditional branch if

(1) if ~ elseif ~ else ~ endif

[Description format]

```
[ Δ ] if [ Δ ] ( Conditional expression 1 ) [ Δ ] [ ( Register name ) ]
    if block
[ Δ ] elseif [ Δ ] ( Conditional expression 2 ) [ Δ ] [ ( Register name ) ]
    elseif block
[ Δ ] else
    else block
[ Δ ] endif
```

[Function]

- if ~ endif

The if block is executed if conditional expression 1 is true.

The if block may occupy several lines.

- if ~ else ~ endif

The if block is executed if conditional expression 1 is true and the else block is executed if it is false.

The if block and else block may occupy several lines.

- if ~ elseif ~ else ~ endif

Several elseif blocks can be written for a single if statement.

If conditional expression 1 is true, the if block is executed. If it is false, conditional expression 2 is tested.

If conditional expression 2 is true, the elseif block is executed. If it is false, the condition of any other elseif that exists prior to the next endif is tested. If there is no elseif, the else block is executed.

The if block, elseif block, and else block may occupy several lines.

[Description]

- Comparison expressions, logic expressions, and test bit expressions can be entered in conditional expressions. If a register name is specified, the specified register is used when testing conditions.

For details of comparison expressions and logic expressions, see "[3.5 Conditional Expressions](#)".

- if ~ else ~ endif is used when coding two branches for a condition.
- if ~ elseif ~ else ~ endif is used when coding several branches for a certain range of values. This differs from a switch statement in that the statement contains a range of values.
- elseif statements and else statements can be omitted and several elseif statements can be entered.

[Generated instructions]

- (1) Processing of if (conditional expression)
 - Generates an instruction to test the condition of the conditional expression.
 - Generates a branch instruction to branch to an elseif block or else block if the condition is not met.
- (2) Processing of elseif (conditional expression)
 - Generates a branch instruction to an endif statement.
 - Generates a label for the branch instruction generated by an if statement.
 - Generates an instruction to test the condition of the conditional expression.
 - Generates a branch instruction to branch to an elseif block or else block if the condition is not met.
- (3) Processing of else
 - Generates a branch instruction to an endif statement.
 - Generates a label for the branch instruction generated by an if statement or elseif statement.
- (4) Processing of endif
 - Generates a label for the branch instruction generated by an if statement, elseif statement, or else statement.
- (5) Additional description
 - These blocks can be mixed with elseif_bit.

[Use examples]

- (1) When entered in lower case letters

Output source	Input source
<pre> CMP A , #0 BNZ \$?L1 BF TFLG.0 , \$?L2 SET1 CY BR ?L3 ?L2 : CLR1 CY ?L3 : MOVW AX , #0FFH BR ?L4 ?L1 : MOVW BC , #0A00H ?L4 :</pre>	<pre> if (A == #0) CY = TFLG.0 else BC = #0A00H endif</pre>

(2) When entered in upper case letters

Output source		Input source
	CMP A , #0	IF (A == #0)
	BZ \$?L5	
	BR ?L6	
?L5 :		
	BF TFLG.0 , \$?L7	CY = TFLG.0
	SET1 CY	
	BR ?L8	
?L7 :		
	CLR1 CY	
?L8 :		
	MOVW AX , #0FFH	AX = #0FFH
	BR ?L9	
?L6 :		ELSE
	MOVW BC , #0A00H	BC = #0A00H
?L9 :		ENDIF

Conditional branch `if_bit`

(2) `if_bit ~ elseif_bit ~ else ~ endif`

[Description format]

```
[ Δ ] if_bit [ Δ ] ( test bit expression 1 )
    if_bit block
[ Δ ] elseif_bit [ Δ ] ( test bit expression 2 )
    elseif_bit block
[ Δ ] else [ Δ ]
    else block
[ Δ ] endif [ Δ ]
```

[Function]

- `if_bit ~ endif`
If conditional expression 1 is true, the `if_bit` block is executed.
The `if_bit` block may occupy several lines.
- `if_bit ~ else ~ endif`
The `if_bit` block is executed if conditional expression 1 is true and the `else` block is executed if it is false.
The `if_bit` block and `else` block may occupy several lines.
- `if_bit ~ elseif_bit ~ else ~ endif`
If conditional expression 1 is true, the `if_bit` block is executed. If it is false, conditional expression 2 is tested.
If conditional expression 2 is true, the `elseif_bit` block is executed. If it is false, the condition of any `elseif_bit` that exists before the next `endif` is tested.
If there is no `elseif_bit`, the `else` block is executed.
The `if_bit` block, `elseif_bit` block, and `else` block may occupy several lines.
- Additional description
These blocks can be mixed with `elseif`.

[Description]

- Test bit expressions are entered as conditional expressions 1 and 2.
For details of test bit expressions, see "[3.5 Conditional Expressions](#)".
- `if_bit ~ else ~ endif` is used when coding two branches for a condition.
`if_bit ~ elseif_bit ~ else ~ endif` is used when checking several bit symbols for multiple branches.
- `elseif_bit` statements and `else` statements can be omitted and several `elseif_bit` statements can be entered.

[Generated instructions]

- (1) Processing of if_bit (bit condition)
 - Generates a true/false instruction for a bit condition.
- (2) Processing of elseif_bit (bit condition)
 - Generates a branch instruction to an endif statement.
 - Generates a label for the branch instruction generated by an if_bit statement.
 - Generates a true/false instruction for a bit condition.
- (3) Processing of else
 - Generates a branch instruction to an endif statement.
 - Generates a label for the branch instruction generated by an if_bit statement or elseif_bit statement.
- (4) Processing of endif
 - Generates a label for the branch instruction generated by an if_bit statement, elseif_bit statement, or else statement.

[Use examples]

- (1) When entered in lower case letters

Output source	Input source
<pre> BT TRFG.0 , \$?L1 SET1 PRTYFLG.3 BR ?L2 ?L1 : BF PGF.0 , \$?L3 MOVW BC , #0FFH BR ?L2 ?L3 : MOV A , # (FG SHR 6) MOV H , A BF PGF.0 , \$?L4 SET1 CY BR ?L5 ?L4 : CLR1 CY ?L5 : CLR1 BUSYFG.2 ?L2 :</pre>	<pre> if_bit (!TRFG.0) PRTYFLG.3=1 elseif_bit (PGF.0) BC = #0FFH else H = # (FG SHR 6) (A) CY = PFG.0 BUSYFG.2 = 0 endif</pre>

(2) When entered in upper case letters

Output source	Input source
<pre> BF TRFG.0 , \$?L6 BR ?L7 ?L6 : SET1 PRTYFLG.3 BR ?L8 ?L7 : BT PGF.0 , \$?L9 BR ?L10 ?L9 : MOVW BC , #0FFH BR ?L8 ?L10 : MOV A , # (FG SHR 6) MOV H , A BF PFG.0 , \$?L11 SET1 CY BR ?L12 ?L11 : CLR1 CY ?L12 : CLR1 BUSYFG.2 ?L8 : </pre>	<pre> IF_BIT (!TRFG.0) PRTYFLG.3 = 1 ELSEIF_BIT (PGF.0) BC = #0FFH ELSE H = # (FG SHR 6) (A) CY = PFG.0 BUSYFG.2 = 0 ENDIF </pre>

Conditional branch switch

(3) switch ~ case ~ default ~ ends

[Description format]

```
[ Δ ] switch [ Δ ] ( [ Δ ] case symbol [ Δ ] ) [ Δ ] [ ( specified register ) ]
[ Δ ] case [ Δ ] Constant :
    Statement_1
[ [ Δ ] case [ Δ ] Constant :
    Statement_2 ]
[ Δ ] [ default : ]
    Statement_N
[ Δ ] ends
```

[Function]

- If the value of the case symbol matches the case constant, the specified statement is executed.
- If the value of the case symbol does not match any case constant and a default statement has been entered, the default statement is executed.
- Normally, a break statement must be entered to skip a switch block.

[Description]

- The possible specifications for "case symbol" depend on the assembly language of the target device.
- If a break statement is not entered, a comparison instruction is executed for the next case statement.
- Constants can be expressed as binary, octal, decimal, hexadecimal, or character string constants. However, since the ST78K0S recognizes constants as character strings, be careful to use only constants that the assembler can recognize as such.
- The case symbol is transferred to the specified register only when a register specification has been made.

[Generated instructions]

(1) Processing of switch statement

- (a) If a register has not been specified, the case symbol is tested and, when necessary, a transfer instruction to A or AX is generated.
- (b) If a register has been specified, the case symbol is transferred to the specified register. However, an error occurs if a comparison instruction cannot be generated. For details, see [Table 3-2](#).

(2) Processing of case statement

- (a) Labels are generated from branch processing from other case statements.
- (b) CMP or CMPW is generated, and if the specified constant does not match, a branch instruction for another case statement, default statement, or ends statement is generated.
 - ?LTRUE : Branch destination label when specified constant matches
 - ?LFALSE : Branch destination label when specified constant does not match

- If the case statement is expressed in lower case letters and a register specification has not been made in the switch statement

```
CMP ( W )      case symbol , #case constant
BNZ           $?LFALSE
```

- If the case statement is expressed in lower case letters and a register specification has been made in the switch statement

```
CAMP ( W )     specified register , #case constant
BNZ           $?LFALSE
```

- If the case statement is expressed in upper case letters and a register specification has not been made in the switch statement

```
          CMAP ( W )      case symbol , #case constant
          BZ             $?LTRUE
          BR             ?LFALSE
?LTRUE :
```

- If the case statement is expressed in upper case letters and a register specification has been made in the switch statement

```
          CAMP ( W )     specified register , #case constant
          BZ             $?LTRUE
          BR             ?LFALSE
?LTRUE :
```

(3) Processing of default statement

Generates a label for the branch instruction from the case statement

(4) Processing of ends statement

Generates a label for the branch instruction from the case statement or break statement

Table 3-2 Generated Instructions for switch Statements

Case symbol		Without register specification	With register specification														
			a	b	f	g	h	i	j	k	n	o	p	q	r		
a	CY																
b	Bit symbol																
c	[HL].β																
d	Byte user symbol	*3		*1							*2						
e	Byte data	*3		*1													
f	A	*3															
g	Byte register	*1		*1													
h	R0	*1		*1													
i	R1																
j	sfr	*1		*1													
k	PSW	*1		*1													
l	Word user symbol			*1							*2						
m	Word data	*2									*2						
n	AX	*3															
o	Word register	*2									*2						
p	RP0																
q	sfrp																
r	SP	*2									*2						
s	Direct access symbol	*1		*1													
t	Indirect access symbol	*1		*1													
u	[DE]	*1		*1													
v	Immediate symbol	*1		*1							*2						

*1 : Generates MOV instruction

*2 : Generates MOVW instruction

*3 : Does not generate transfer instruction

Empty columns indicate errors.

[Use examples]

(1) When entered in lower case letters

Output source	Input source
<pre> MOV A , R0 CMP A , #1 BNZ \$?L1 BF P1.0 , \$?L2 BTM.3 ?L2 : BR ?L3 ?L1 : CMP A , #2 BNZ \$?L4 BR ?L3 ?L4 : CMP A , #3 BNZ \$?L5 BR ?L3 ?L5 : ?L3 :</pre>	<pre> SWITCH (R0) case 1 : if_bit (P1.0) BTM.3 endif break case 2 : break case 3 : break default : ENDS</pre>

(2) When entered in upper case letters

Output source	Input source
<pre> MOV A , R0 CMP A , #1 BZ \$?L6 BR ?L7 ?L6 : BF P1.0 , \$?L8 BTM.3 ?L8 : BR ?L9 ?L7 : CMP A , #2 BZ \$?L10 BR ?L11 ?L10 : BR ?L9 ?L11 : CMP A , #3 BZ \$?L12 BR ?L13 ?L12 : BR ?L9 ?L13 : ?L9 :</pre>	<pre> SWITCH (R0) CASE 1 : if_bit (P1.0) BTM.3 endif break CASE 2 : break CASE 3 : break DEFAULT : ENDS</pre>

3.4.2 Conditional loop

Conditional loop for

(4) for ~ next

[Description format]

```
[ Δ ] for [ Δ ] ( [ expression 1 ] ; [ expression 2 ] ; [ expression 3 ] ) [
Δ ] [ ( register specification ) ]
    Instruction group
[ Δ ] next
```

[Function]

- The initial value is set by expression 1 and the statement and expression 3 are executed as long as the conditional expression in expression 2 is met. Usually, expression 3 is an [Increment \(++ \)](#) or [Decrement \(- - \)](#) operation.

The meaning is similar to the example shown below.

```
Expression 1
while ( expression 2 )
    Instruction group
    Expression 3
endw
```

[Description]

- (1) Be sure to note that the similar example shown above does not apply to generated instructions.
- (2) The following are entered in expression 1, expression 2, and expression 3.
 - Expression 1 : Initial value setting (assignment expression)
 - Expression 2 : Conditional expression
 - Expression 3 : Increment or decrement expression
- (3) Assignment operators and exchange statements can be entered in expression 1 or expression 3, but when doing so, the conversion output should be checked and modified if necessary.
- (4) It is possible to omit expression 1, expression 2, or expression 3. However, if expression 2 is omitted, an endless loop will occur.
- (5) "forever" can be entered in a conditional expression.
- (6) Since expression 2 and expression 3 control for ~ next, the contents of these expressions should not be changed by an executable statement. Changing these contents can result in faulty operation.

[Generated instructions]

- (1) Processing of for statement (expression 1 ; expression 2 ; expression 3)
 - (a) Generates instruction for expression 1. If a register name has been specified, the specified register is used for assignments and comparisons.
 - (b) Generates a branch instruction to the statement that tests expression 2's conditions.
 - (c) Generates a label for the branch instruction generated by a next statement.
 - (d) Generates a label for the branch instruction generated at (b).
 - (e) Generates a condition testing instruction for expression 2.
- (2) Processing of next statement
 - (a) Generates a branch instruction to the label generated via for statement processing (c).
 - (b) Generates a label for the branch instruction for skipping a for block.
 - (c) Generates an instruction for expression 3's assignment expression.
- (3) Additional description

The following method is recommended for more effective use of for ~ next statements.

- Use saddr instead of a register name as the control variable in expression 1 and expression 3.
- When specifying a register, specify either A or AX.
- When executing a loop for at least 256 repetitions, nest a for statement and use two saddr variables as the control variables.

Remark The above method is recommended because of the limited range of symbols that can be entered as operands in order to output CMP or CMPW as generated instructions for the conditional expression in expression 2.

[Use examples]

(1) When entered in lower case letters

Output source	Input source
<pre>?L1 : MOV i , #0H CMP i , #0FFH BNC \$?L2 CALL !XXX NC i BR ?L1 ?L2 :</pre>	<pre>for (i = #0H ; i < #0FFH ; i++) CALL !XXX next</pre>

(2) When entered in upper case letters

Output source	Input source
<pre>?L3 : MOV i , #0H CMP i , #0FFH BC \$?L4 BR ?L5 ?L4 : CALL !XXX INC i BR ?L3 ?L5 :</pre>	<pre>FOR (i = #0H ; i < #0FFH ; i++) CALL !XXX NEXT</pre>

Conditional loop while

(5) while ~ endw

[Description format]

```
[ Δ ] while [ Δ ] ( conditional expression ) [ Δ ] [ ( register specification ) ]
        Instruction group
[ Δ ] endw
```

[Function]

- The instruction group is repeatedly executed as long as the conditional expression remains true.

[Description]

- It is possible to enter comparison expressions, logic expressions, test bit expressions, and "forever" as conditional expressions.
If "forever" is entered, the result is an endless loop.
- As the register name, specify the register used in the comparison expression or logic expression entered as "(conditional expression)".
- Since the conditional expression is tested before the instruction group is executed, if the first conditional expression is found to be false, the instruction group is not executed even once.

[Generated instructions]

- (1) Processing of while (conditional expression) statement
 - Generates a label for the branch instruction generated by endw.
 - Generates a condition testing instruction. If a register name has been specified, the specified register is used when generating the condition testing instruction.
 - Generates a branch instruction for removing the while (conditional expression) statement from the while block when the condition tests as false.
- (2) endw
 - Generates a branch instruction for an execution loop.
 - Generates a label for the branch instruction that is used to remove endw from the while block.

[Use examples]

(1) When entered in lower case letters

Output source	Input source
<pre>?L1 : CMPW AX , #0FFFH BNC \$?L2 MOV B , #0FH INCW HL BR ?L1 ?L2 :</pre>	<pre>while (AX < #0FFFH) B = #0FH HL++ endw</pre>

(2) When entered in upper case letters

Output source	Input source
<pre>?L3 : CMPW AX , #0FFFH BC \$?L4 ?L4 : MOV B , #0FH NCW HL BR ?L3 ?L5 :</pre>	<pre>WHILE (AX < #0FFFH) B = #0FH HL++ ENDW</pre>

Conditional loop `while_bit`

(6) `while_bit ~ endw`

[Description format]

```
[ Δ ] while_bit [ Δ ] ( bit condition )
        Instruction group
[ Δ ] endw
```

[Function]

- The instruction group can be executed as long as the bit condition is true.

[Description]

- Since the bit condition is tested before the instruction group is executed, if the first bit condition is found to be false, the instruction group is not executed even once.

[Generated instructions]

- (1) Processing of `while_bit (bit condition)` statement
 - Generates a label for the branch instruction generated by `endw`.
 - Generates an instruction for testing the bit condition as true or false.
 - Generates a branch instruction for removing the `while_bit` statement from the `while_bit ~ endw` block when the bit condition tests as false.
- (2) Processing of `endw`
 - Generates a branch instruction for an execution loop.
 - Generates a label for the branch instruction that is used to remove `endw` from the `while_bit` block.

[Use examples]

(1) When entered in lower case letters

Output source	Input source
<pre>?L1 : BT TRFG.0 , \$?L2 MOV A , PORT1 CMP A , #04H BNZ \$?L3 MOV X , #0FFH BR ?L4 ?L3 : CLR1 PFG.0 ?L4 : BR ?L1 ?L2 :</pre>	<pre>while_bit (!TRFG.0) A = PORT1 if (A == #04H) X = #0FFH else PFG.0 = 0 endif endw</pre>

(2) When entered in upper case letters

Output source	Input source
<pre>?L5 : BF TRFG.0 , \$?L6 BR ?L7 ?L6 : MOV A , PORT1 CMP A , #04H BNZ \$?L8 MOV X , #0FFH BR ?L9 ?L8 : CLR1 PFG.0 ?L9 : BR ?L5 ?L7 :</pre>	<pre>WHILE_BIT (!TRFG.0) A = PORT1 if (A == #04H) X = #0FFH else PFG.0 = 0 endif ENDW</pre>

Conditional loop until

(7) repeat ~ until

[Description format]

```
[ Δ ] repeat
      Instruction group
[ Δ ] until [ Δ ] ( conditional expression ) [ Δ ] [ ( register specification
) ]
```

[Function]

- The instruction group is repeatedly executed as long as the conditional expression remains true.

[Description]

- It is possible to enter comparison expressions, logic expressions, test bit expressions, and "forever" as conditional expressions.
If "forever" is entered, the result is an endless loop.
- As the register name, specify the register used in the comparison expression or logic expression entered as "(conditional expression)".
- The conditional expression is tested after the instruction group is executed. Therefore, if the first conditional expression is found to be true, the instruction group is executed once.

[Generated instructions]

- (1) Processing of repeat statement
 - Generates a label for the branch instruction generated by until.
- (2) Processing of until (conditional expression) statement
 - Generates a condition testing instruction for the conditional expression.
 - Generates a branch instruction for the label that was generated by repeat in order to execution the instruction group during repeat ~ until and while the conditional expression tests as false. If the conditional expression tests as true, the until statement is removed from the repeat block.

[Use examples]

(1) When entered in lower case letters

Output source	Input source
<pre>?L1 : MOVW AX , BC CMP ABC , #0CH BNZ \$?L2 CALL !XXX ?L2 : INC CNT CMP CNT , #0FFH BNZ \$?L1</pre>	<pre>repeat AX = BC if (ABC == #0CH) CALL !XXX endif CNT++ until (CNT == #0FFH)</pre>

(2) When entered in upper case letters

Output source	Input source
<pre>?L3 : MOVW AX , BC CMP ABC , #0CH BNZ \$?L4 CALL !XXX ?L4 : INC CNT CMP CNT , #0FFH BZ \$?L5 BR ?L3 ?L5 :</pre>	<pre>REPEAT AX = BC if (ABC == #0CH) CALL !XXX endif CNT++ UNTIL (CNT == #0FFH)</pre>

Conditional loop until_bit

(8) repeat ~ until_bit

[Description format]

```
[ Δ ] repeat
      Instruction group
[ Δ ] until_bit [ Δ ] ( test bit expression )
```

[Function]

- The instruction group is repeatedly executed as long as the bit condition is false.

[Description]

- The bit condition is tested after the instruction group is executed. Therefore, if the first bit condition is found to be true, the instruction group is executed once.

[Generated instructions]

- (1) Processing of repeat
 - Generates a label for the branch instruction generated by until_bit.
- (2) Processing of until_bit (bit condition)
 - Generates a branch instruction for the label that is generated by repeat in order to execute the instruction group between repeat and until_bit when the conditional expression tests as false. If the conditional expression tests as true, until_bit is removed from the repeat block.

[Use examples]

- (1) When entered in lower case letters

Output source	Input source
<pre>?L1 : MOV B , #8H CALL !XXX BF TRFG.0 , \$?L1</pre>	<pre>repeat B = #8H CALL !XXX until_bit (TRFG.0)</pre>

- (2) When entered in upper case letters

Output source	Input source
<pre>?L2 : MOV B , #8H CALL !XXX BT TRFG.0 , \$?L3 BR ?L2 ?L3 :</pre>	<pre>REPEAT B = #8H CALL !XXX UNTIL_BIT (TRFG.0)</pre>

Conditional loop break

(9) break

[Description format]

```
[ Δ ] break
```

[Function]

- Terminates execution of the innermost nested block among while, repeat, for, and switch blocks.

[Description]

- An error occurs if a statement other than a while, while_bit, repeat ~ until, repeat ~ until_bit, for, or switch statement has been entered.

[Generated instructions]

- Generates an unconditional branch instruction to remove while, repeat, for, or switch blocks.

```
BR      ?Lxxxx
```

[Use example]

Output source	Input source
<pre>?L1 : MOV X , #0 MOV PORT4 , A CMP A , #0FH BNZ \$?L2 BR ?L3 ?L2 : INCW HL BR ?L1 ?L3 :</pre>	<pre>while (forever) X = #0 PORT4 = A if (A == #0FH) break endif HL++ endw</pre>

Conditional loop continue

(10) continue

[Description format]

```
[ Δ ] continue
```

[Function]

- Skips processing following continue within the innermost nested block among a while, while_bit, repeat ~ until, repeat ~ until_bit, or for statement and sets an unconditional branch before the condition is tested.

[Description]

- This is used to skip subsequent processing from the middle of a block and execute the next loop.
- An error occurs if a statement other than a while, while_bit, repeat ~ until, repeat ~ until_bit, or for statement has been entered.

[Generated instructions]

- Generates an unconditional branch instruction for a label to repeat a while, while_bit, repeat ~ until, repeat ~ until_bit, or for block

```
BR      ?Lxxxx
```

[Use example]

Output source	Input source
<pre>?L1 : CMP SYM , #0FH BNZ \$?L2 MOV B , #0 MOV PORT4 , A CMP A , #0FH BNZ \$?L3 BR ?L1 BR ?L4 ?L3 : INCW HL ?L4 : BR ?L1 ?L2 :</pre>	<pre>while (SYM == #0FH) B = #0 PORT4 = A if (A == #0FH) continue else HL++ endif endw</pre>

Conditional loop goto

(11) goto

[Description format]

```
[ Δ ] goto Δ label
```

[Function]

- Unconditionally branches to a label.

[Description]

- goto statements are entered when immediate error processing is required such as in an error processing program, or when collective processing of errors at multiple locations is needed.
- The symbols shown in the assembly language label column are specified as label names.

[Generated instructions]

- (1) Generates the following instruction.

```
BR      Label
```

- (2) The goto statement's labels are not automatically generated by the ST78K0S. Note also that the ST78K0S does not automatically check whether or not a branch destination label exists.

[Use examples]

Output source	Input source
<pre>?L1 : MOV B , #0 MOV PORT4 , A CMP A , #0FH BNZ \$?L2 BR ERROR ?L2 : INCW HL BR ?L1</pre>	<pre>while (forever) B = #0 PORT4 = A if (A == #0FH) goto ERROR endif HL++ endw</pre>

3.5 Conditional Expressions

Conditional expressions are used to set conditions via control statements.

The following are examples of conditional expressions.

- **Comparison expressions** : Compares first and second values and tests them as true or false.
- **Test bit expressions** : Determines flag on/off status based on bit symbols.
- **Logical operations** : Performs a logical operation for a conditional expression when conditions are combined.

If (γ) is specified at the end of a comparison, a comparison can be made between α and β values that cannot be compared directly.

γ specifies the register that is used for this comparison.

3.5.1 Comparison expressions

In the description of each comparison expression, "?LTRUE" is used as the branch destination label for when the comparison tests as true and "?LFALSE" is used as the branch destination label when it tests as false.

See "3.3 Register Specification" for a description of the register specification Description format.

The ST78K0S does not test whether or not the symbols entered on the left and right sides of a comparison expression are entered correctly as assembly language operands. However, a data size test is performed, as described in "2.6 Data Sizes" to determine whether or not an instruction can be generated. In addition, when specifying a register, the possibility of generating an instruction using the specified register is tested.

An error message is output when a test results in an error.

For details, see the relevant generated instruction.

The following pages describe the functions of the various comparison expressions.

The use examples show as comment statements the source files to which generated instructions are input.

Table 3-3 Generated Instructions for Comparison Instructions

Symbol			β																					
			a	b	c	d	e	f	g	h	i	j	k	l	m	n	o	p	q	r	s	t	u	v
α	a	CY																						
	b	Bit symbol																						
	c	[HL]. β																						
	d	Byte user symbol																						*1
	e	Byte data																						*1
	f	A			*1	*1		*1	*1	*1			*1								*1	*1		*1
	g	Byte register																						
	h	R0																						
	i	R1																						
	j	sfr																						
	k	PSW																						
	l	Word user symbol																						
	m	Word data																						*2
	n	AX																						
	o	BC, DE, HL																						
	p	RP0, RP1, RP2, RP3																						
	q	sfrp																						
	r	SP																						
	s	Direct access symbol																						
	t	Indirect access symbol																						
u	[DE]																							
v	Immediate symbol																							

*1 : Generates CMP instruction

*2 : Generates CMPW instruction

Empty columns indicate errors.

Table 3-4 Comparison Expressions

Comparison Expression	Description Format	Function
Equal (==)	$\alpha == \beta$	True when $\alpha = \beta$, false when $\alpha \neq \beta$
NotEqual (!=)	$\alpha != \beta$	True when $\alpha \neq \beta$, false when $\alpha = \beta$
LessThan (<)	$\alpha < \beta$	True when $\alpha < \beta$, false when $\alpha \geq \beta$
GreaterThan (>)	$\alpha > \beta$	True when $\alpha > \beta$, false when $\alpha \leq \beta$
GreaterEqual (>=)	$\alpha \geq \beta$	True when $\alpha \geq \beta$, false when $\alpha < \beta$
LessEqual (<=)	$\alpha \leq \beta$	True when $\alpha \leq \beta$, false when $\alpha > \beta$
FOREVER (forever)	<code>forever</code>	Endlessly loops the loop statement

Comparison expressions Equal (==)

(1) Equal (==)

[Description format]

```
[ Δ ] [ size specification ] α [ Δ ] == [ Δ ] [ size specification ] [ Δ ] β
[ Δ ] [ ( register specification ) ]
```

[Function]

- When there is no register specification
True when the contents of α and β are equal, false when they are not equal.
- When there is a register specification
The contents of α are transferred to the specified register. True is the result when the contents of the specified register are equal to the contents of β and false is the result when they are not equal.

[Description]

- When there is no register specification
For α and β , be sure to specify contents that can be entered in CMP or CMPW.
- When there is a register specification
For α , be sure to specify contents that can be entered in MOV or MOVW.
For β , be sure to specify contents that can be entered in CMP or CMPW.

[Generated instructions]

- (1) If the control statement is entered in lower case letters and there is no register specification

```
CMP ( W )      α , β
BNZ           $?LFALSE
```

- (2) If the control statement is entered in lower case letters and there is a register specification

```
MOV ( W )      Specified register , α
CMP ( W )      Specified register , β
BNZ           $?LFALSE
```

- (3) If the control statement is entered in upper case letters and there is no register specification

```
      CMP ( W )      α , β
      BZ             $?LTRUE
      BR             ?LFALSE
?LTRUE :
```

- (4) If the control statement is entered in upper case letters and there is a register specification

MOV (W)	Specified register , α
CMP (W)	Specified register , β
BZ	$\$?LTRUE$
BR	$?LFALSE$
$?LTRUE$:	

For details of combinations of α and β , see [Table 3-3](#). α indicates the specified register. For further description of generated instructions for MOV, see "4.2 (1) Assign (=)".

[Use examples]

- (1) If the control statement is entered in lower case letters and there is no register specification

Output source	Input source
<pre> CMPW AX , #0F0FH BNZ \$?L1 CALL !XXX BR ?L2 ?L1 : CALL !YYY ?L2 :</pre>	<pre> if (AX == #0F0FH) CALL !XX else CALL !YYY endif</pre>

- (2) If the control statement is entered in lower case letters and there is a register specification

Output source	Input source
<pre> MOV A , !XYZ CMP A , #5 BNZ \$?L3 CALL !PPP ?L3 :</pre>	<pre> if (!XYZ == #5 (A)) CALL !PPP endif</pre>

- (3) If the control statement is entered in upper case letters and there is no register specification

Output source	Input source
<pre> CMPW AX , #0F0FH BZ \$?L4 BR ?L5 ?L4 : CALL !XXX BR ?L6 ?L5 : CALL !YYY ?L6 :</pre>	<pre> IF (AX == #0F0FH) CALL !XXX ELSE CALL !YYY ENDIF</pre>

- (4) If the control statement is entered in upper case letters and there is a register specification

Output source	Input source
<pre> MOV A , !XYZ CMP A , #5 BZ \$?L7 BR ?L8 ?L7 : CALL !PPP ?L8 :</pre>	<pre> IF (!XYZ == #5 (A)) CALL !PPP ENDIF</pre>

Comparison expressions NotEqual (!=)

(2) NotEqual (!=)

[Description format]

```
[ Δ ] [ size specification ] [ Δ ] α [ Δ ] != [ Δ ] [ size specification ] [ Δ ] β [ Δ ] [ ( register specification ) ]
```

[Function]

- When there is no register specification
True when the contents of α and β are not equal, false when they are equal.
- When there is a register specification
The contents of α are transferred to the specified register. True is the result when the contents of the specified register are not equal to the contents of β and false is the result when they are equal.

[Description]

- When there is no register specification
For α and β , be sure to specify contents that can be entered in CMP or CMPW.
- When there is a register specification
For α , specify contents that can be entered in MOV or MOVW.
For β , specify contents that can be entered in CMP or CMPW.

[Generated instructions]

- (1) If the control statement is entered in lower case letters and there is no register specification

```
CMP ( W )      α , β
BZ             $?LFALSE
```

- (2) If the control statement is entered in lower case letters and there is a register specification

```
MOV ( W )      Specified register , α
CMP ( W )      Specified register , β
BZ             $?LFALSE
```

- (3) If the control statement is entered in upper case letters and there is no register specification

```
      CMP ( W )      α , β
      BNZ            $?LTRUE
      BR             ?LFALSE
?LTRUE :
```

- (4) If the control statement is entered in upper case letters and there is a register specification

MOV (W)	Specified register , α
CMP (W)	Specified register , β
BNZ	$\$?LTRUE$
BR	$?LFALSE$
$?LTRUE$:	

For details of combinations of α and β , see [Table 3-3](#). α indicates the specified register. For further description of generated instructions for MOV, see "4.2 (1) Assign (=)".

[Use examples]

- (1) If the control statement is entered in lower case letters and there is no register specification

Output source	Input source
<pre> CMPW AX , #0FFFH BZ \$?L1 CALL !XXX BR ?L2 ?L1 : CALL !YYY ?L2 :</pre>	<pre> if (AX != #0FFFH) CALL !XXX else CALL !YYY endif</pre>

- (2) If the control statement is entered in lower case letters and there is a register specification

Output source	Input source
<pre> MOV A , !XYZ CMP A , #5 BZ \$?L CALL !PPP ?L3 :</pre>	<pre> if (!XYZ != #5 (A)) CALL !PPP endif</pre>

- (3) If the control statement is entered in upper case letters and there is no register specification

Output source	Input source
<pre> CMPW AX , #0FFFH BNZ \$?L4 BR ?L5 ?L4 : CALL !XXX BR ?L6 ?L5 : CALL !YYY ?L6 :</pre>	<pre> IF (AX != #0FFFH) CALL !XXX ELSE CALL !YYY ENDIF</pre>

- (4) If the control statement is entered in upper case letters and there is a register specification

Output source	Input source
<pre> MOV A , !XYZ CMP A , #5 BNZ \$?L7 BR ?L8 ?L7 : CALL !PPP ?L8 :</pre>	<pre> IF (!XYZ != #5 (A)) CALL !PPP ENDIF</pre>

Comparison expressions LessThan (<)

(3) LessThan (<)

[Description format]

```
[ Δ ] [ size specification ] [ Δ ] α [ Δ ] < [ Δ ] [ size specification ] [ Δ ] β [ Δ ] [ ( register specification ) ]
```

[Function]

- When there is no register specification
True when the contents of α are less than the contents of β , false when otherwise (i.e., equal to or greater than).
- When there is a register specification
The contents of α are transferred to the specified register. True is the result when the contents of the specified register are less than the contents of β and false is the result when they are otherwise.

[Description]

- When there is no register specification
For α and β , be sure to specify contents that can be entered in CMP or CMPW.
- When there is a register specification
For α , be sure to specify contents that can be entered in MOV or MOVW.
For β , be sure to specify contents that can be entered in CMP or CMPW.

[Generated instructions]

- (1) If the control statement is entered in lower case letters and there is no register specification

```
CMP ( W )      α , β
BNC            $?LFALSE
```

- (2) If the control statement is entered in lower case letters and there is a register specification

```
MOV ( W )      Specified register , α
CMP ( W )      Specified register , β
BNC            $?LFALSE
```

- (3) If the control statement is entered in upper case letters and there is no register specification

```
      CMP ( W )      α , β
      BC             $?LTRUE
      BR             ?LFALSE
?LTRUE :
```

- (4) If the control statement is entered in upper case letters and there is a register specification

MOV (W)	Specified register , α
CMP (W)	Specified register , β
BC	$\$?LTRUE$
BR	$?LFALSE$
$?LTRUE$:	

For details of combinations of α and β , see [Table 3-3](#). α indicates the specified register. For further description of generated instructions for MOV, see "4.2 (1) Assign (=)".

[Use examples]

- (1) If the control statement is entered in lower case letters and there is no register specification

Output source	Input source
CMP A , [HL]	if (A < [HL])
BNC $\$?L1$	
CALL !XXX	CALL !XXX
BR $?L2$	
$?L1$:	else
CALL !YYY	CALL !YYY
$?L2$:	endif

- (2) If the control statement is entered in lower case letters and there is a register specification

Output source	Input source
MOVW AX , ABCP	if (ABCP < #0FE00H (AX))
CMPW AX , #0FE00H	
BNC $\$?L3$	
CALL !PPP	CALL !PPP
$?L3$:	endif

- (3) If the control statement is entered in upper case letters and there is no register specification

Output source	Input source
CMP A , [HL]	IF (A < [HL])
BC $\$?L4$	
BR $?L5$	
$?L4$:	CALL !XXX
CALL !XXX	
BR $?L6$	
$?L5$:	ELSE
CALL !YYY	CALL !YYY
$?L6$:	ENDIF

- (4) If the control statement is entered in upper case letters and there is a register specification

Output source	Input source
<pre> MOVW AX , ABCP CMPW AX , #0FE00H BC \$?L7 BR ?L8 ?L7 : CALL !PPP ?L8 : </pre>	<pre> IF (ABCP < #0FE00H (AX)) CALL !PPP ENDIF </pre>

Comparison expressions GreaterThan (>)

(4) GreaterThan (>)

[Description format]

```
[ Δ ] [ size specification ] [ Δ ] α [ Δ ] > [ Δ ] [ size specification ] [ Δ ]
] β [ Δ ] [ ( register specification ) ]
```

[Function]

- When there is no register specification
True when the contents of α are greater than the contents of β , false when otherwise (i.e. equal to or less than).
- When there is a register specification
The contents of α are transferred to the specified register. True is the result when the contents of the specified register are greater than the contents of β and false is the result when they are otherwise.

[Description]

- When there is no register specification
For α and β , be sure to specify contents that can be entered in CMP or CMPW.
- When there is a register specification
For α , be sure to specify contents that can be entered in MOV or MOVW.
For β , be sure to specify contents that can be entered in CMP or CMPW.

[Generated instructions]

- (1) If the control statement is entered in lower case letters and there is no register specification

```
CMP ( W )      α , β
BZ             $?LFALSE
BC             $?LFALSE
```

- (2) If the control statement is entered in lower case letters and there is a register specification

```
MOV ( W )      Specified register , α
CMP ( W )      Specified register , β
BZ             $?LFALSE
BC             $?LFALSE
```

- (3) If the control statement is entered in upper case letters and there is no register specification

```
      CMP ( W )      α , β
      BZ             $$ + 4
      BNC            $?LTRUE
      BR             ?LFALSE
?LTRUE :
```

- (4) If the control statement is entered in upper case letters and there is a register specification

MOV (W)	Specified register , α
CMP (W)	Specified register , β
BZ	$$$ + 4$
BNC	$?\text{LTRUE}$
BR	$?\text{LFALSE}$
$?\text{LTRUE} :$	

For details of combinations of α and β , see [Table 3-3](#). α indicates the specified register. For further description of generated instructions for MOV, see "4.2 (1) Assign (=)".

[Use examples]

- (1) If the control statement is entered in lower case letters and there is no register specification

Output source	Input source
CMP A , [HL]	if (A > [HL])
BZ $?\text{L1}$	
BC $?\text{L1}$	CALL !XXX
CALL !XXX	
BR $?\text{L2}$	else
$?\text{L1} :$	CALL !YYY
CALL !YYY	endif
$?\text{L2} :$	

- (2) If the control statement is entered in lower case letters and there is a register specification

Output source	Input source
MOVW AX , ABCP	if (ABCP > #0FE40H (AX))
CMPW AX , #0FE40H	
BZ $?\text{L3}$	
BC $?\text{L3}$	CALL !PPP
CALL !PPP	
$?\text{L3} :$	endif

- (3) If the control statement is entered in upper case letters and there is no register specification

Output source	Input source
CMP A , [HL]	IF (A > [HL])
BZ $$$ + 4$	
BNC $?\text{L4}$	
BR $?\text{L5}$	CALL !XXX
$?\text{L4} :$	
CALL !XXX	
BR $?\text{L6}$	ELSE
$?\text{L5} :$	CALL !YYY
CALL !YYY	ENDIF
$?\text{L6} :$	

- (4) If the control statement is entered in upper case letters and there is a register specification

Output source	Input source
<pre> MOVW AX , ABCP CMPW AX , #0FE40H BZ \$\$ + 4 BNC \$?L7 BR ?L8 ?L7 : CALL !PPP ?L8 : </pre>	<pre> IF (ABCP > #0FE40H (AX)) CALL !PPP ENDIF </pre>

Comparison expressions GreaterEqual (>=)

(5) GreaterEqual (>=)

[Description format]

```
[ Δ ] [ size specification ] [ Δ ] α [ Δ ] >= [ Δ ] [ size specification ] [ Δ ] β [ Δ ] [ ( register specification ) ]
```

[Function]

- When there is no register specification
True when the contents of α are greater than or equal to the contents of β , false when they are less than the contents of β .
- When there is a register specification
The contents of α are transferred to the specified register. True is the result when the contents of the specified register are greater than or equal to the contents of β and false is the result when they are less than the contents of β .

[Description]

- When there is no register specification
For α and β , be sure to specify contents that can be entered in CMP or CMPW.
- When there is a register specification
For α , be sure to specify contents that can be entered in MOV or MOVW.
For β , be sure to specify contents that can be entered in CMP or CMPW.

[Generated instructions]

- (1) If the control statement is entered in lower case letters and there is no register specification

```
CMP ( W )      α , β
BC              $?LFALSE
```

- (2) If the control statement is entered in lower case letters and there is a register specification

```
MOV ( W )      Specified register , α
CMP ( W )      Specified register , β
BC              $?LFALSE
```

- (3) If the control statement is entered in upper case letters and there is no register specification

```
      CMP ( W )      α , β
      BNC            $?LTRUE
      BR             ?LFALSE
?LTRUE :
```

- (4) If the control statement is entered in upper case letters and there is a register specification

MOV (W)	Specified register , α
CMP (W)	Specified register , β
BNC	$\$?LTRUE$
BR	$?LFALSE$
$?LTRUE$:	

For details of combinations of α and β , see [Table 3-3](#). α indicates the specified register. For further description of generated instructions for MOV, see "4.2 (1) Assign (=)".

[Use examples]

- (1) If the control statement is entered in lower case letters and there is no register specification

Output source	Input source
<pre> CMP A , [HL] BC \$?L1 CALL !XXX BR ?L2 ?L1 : CALL !YYY ?L2 :</pre>	<pre> if (A >= [HL]) CALL !XXX else CALL !YYY endif</pre>

- (2) If the control statement is entered in lower case letters and there is a register specification

Output source	Input source
<pre> MOVW AX , DE CMPW AX , #0FE30H BC \$?L3 CALL !PPP ?L3 :</pre>	<pre> if (DE >= #0FE30H (AX)) CALL !PPP endif</pre>

- (3) If the control statement is entered in upper case letters and there is no register specification

Output source	Input source
<pre> CMP A , [HL] BNC \$?L4 BR ?L5 ?L4 : CALL !XXX BR ?L6 ?L5 : CALL !YYY ?L6 :</pre>	<pre> IF (A >= [HL]) CALL !XXX ELSE CALL !YYY ENDIF</pre>

- (4) If the control statement is entered in upper case letters and there is a register specification

Output source	Input source
<pre> MOVW AX , DE CMPW AX , #0FE30H BNC \$?L7 BR ?L8 ?L7 : CALL !PPP ?L8 : </pre>	<pre> IF (DE >= #0FE30H (AX)) CALL !PPP ENDIF </pre>

Comparison expressions LessEqual (<=)

(6) LessEqual (<=)

[Description format]

```
[ Δ ] [ size specification ] [ Δ ] α [ Δ ] <= [ Δ ] [ size specification ] [ Δ ] β [ Δ ] [ ( register specification ) ]
```

[Function]

- When there is no register specification
True when the contents of α are less than or equal to the contents of β , false when they are greater than the contents of β .
- When there is a register specification
The contents of α are transferred to the specified register. True is the result when the contents of the specified register are less than or equal to the contents of β and false is the result when they are greater than the contents of β .

[Description]

- When there is no register specification
For α and β , be sure to specify contents that can be entered in CMP or CMPW.
- When there is a register specification
For α , be sure to specify contents that can be entered in MOV or MOVW.
For β , be sure to specify contents that can be entered in CMP or CMPW.

[Generated instructions]

- (1) If the control statement is entered in lower case letters and there is no register specification

```
CMP ( W )      α , β
BZ             $$ + 4
BNC           $?LFALSE
```

- (2) If the control statement is entered in lower case letters and there is a register specification

```
MOV ( W )      Specified register , α
CMP ( W )      Specified register , β
BZ             $$ + 4
BNC           $?LFALSE
```

- (3) If the control statement is entered in upper case letters and there is no register specification

```
      CMP ( W )      α , β
      BZ             $?LTRUE
      BC             $?LTRUE
      BR             ?LFALSE
?LTRUE :
```

- (4) If the control statement is entered in upper case letters and there is a register specification

MOV (W)	Specified register , α
CMP (W)	Specified register , β
BZ	\$?LTRUE
BC	\$\$?LTRUE
BR	?LFALSE
?LTRUE :	

For details of combinations of α and β , see [Table 3-3](#). α indicates the specified register. For further description of generated instructions for MOV, see "[4.2 \(1\) Assign \(= \)](#)".

[Use examples]

- (1) If the control statement is entered in lower case letters and there is no register specification

Output source	Input source
<pre> CMP A , [HL] BZ \$\$ + 4 BNC \$?L1 CALL !XXX BR ?L2 ?L1 : CALL !YYY ?L2 :</pre>	<pre> if (A <= [HL]) CALL !XXX else CALL !YYY endif</pre>

- (2) If the control statement is entered in lower case letters and there is a register specification

Output source	Input source
<pre> MOVW AX , HL CMPW AX , #0FE20H BZ \$\$ + 4 BNC \$?L3 CALL !PPP ?L3 :</pre>	<pre> if (HL <= #0FE20H (AX)) CALL !PPP endif</pre>

- (3) If the control statement is entered in upper case letters and there is no register specification

Output source	Input source
<pre> CMP A , [HL] BZ \$?L4 BC \$?L4 BR ?L5 ?L4 : CALL !XXX BR ?L6 ?L5 : CALL !YYY ?L6 :</pre>	<pre> IF (A <= [HL]) CALL !XXX ELSE CALL !YYY ENDIF</pre>

- (4) If the control statement is entered in upper case letters and there is a register specification

Output source	Input source
<pre> MOVW AX , HL CMPW AX , #0FE20H BZ \$?L7 BC \$?L7 BR ?L8 ?L7 : CALL !PPP ?L8 : </pre>	<pre> IF (HL <= #0FE20H (AX)) CALL !PPP ENDIF </pre>

Comparison expressions FOREVER (forever)

(7) FOREVER (forever)

[Description format]

```
[ Δ ] forever [ Δ ]
```

[Function]

- Sets loop statement as an endless loop, without generating a compare instruction.

[Description]

- Can be entered in a loop statement (for statement, while statement, until statement) type of conditional expression.

[Use examples]

- (1) for statement

Output source	Input source
<pre>?L1 : MOV i , #0 MOV A , i CALL !XXX CMPW AX , #0FFH BNZ \$?L2 BR ?L3 ?L2 : INC i BR ?L1 ?L3 :</pre>	<pre>for (i = #0 ; forever ; i++) A = i CALL !XXX if (AX == #0FFH) break endif next</pre>

- (2) while statement

Output source	Input source
<pre>?L4 : BF forever , \$?L5 MOV A , i CALL !XXX CMPW AX , #0FFH BNZ \$?L6 BR ?L5 ?L6 : INC i BR ?L4 ?L5 :</pre>	<pre>while (forever) A = i CALL !XXX if (AX == #0FFH) break endif i++ endw</pre>

(3) repeat statement

Output source	Input source
<pre>?L7 : MOV A , i CALL !XXX CMPW AX , #0FFH BNZ \$?L8 BR ?L9 ?L8 : INC i BR ?L7 ?L9 :</pre>	<pre>repeat A = i CALL !XXX if (AX == #0FFH) break endif i++ until (forever)</pre>

3.5.2 Test bit expressions

In the description of each type of test bit expression, it is noted that "?LTRUE" is used as the branch destination label when the test result is true and "?LFALSE" is used as this label when the test result is false.

The ST78K0S does not test whether or not test bit expression code is entered correctly as assembly language operands. However, a data size test is performed, as described in "2.6 Data Sizes".

In addition, "Z" is also processed as a bit symbol.

The ST78K0S does not use the assembler's directive (EQU) to check whether or not a bit symbol has been defined. However, user symbols can also be processed as bit symbols.

An error message is output when the test result is an error.

For details, see the particular generating instruction.

The following pages describe the functions of the various test bit expressions.

The use examples show as comment statements the source files to which generated instructions are input.

Table 3-5 Test Bit Expressions

Test Bit Expression	Description Format	Function
Bit symbol	Bit symbol	True when specified bit is 1
!bit symbol	!bit symbol	True when specified bit is 0

Test bit expressions Positive logic (bit)

(1) Bit symbol

[Description format]

```
[ Δ ] bit symbol [ Δ ]
```

[Function]

- True when the bit symbol contents are 1, false when they are 0.
- The following control statements are able to include bit symbols entered as conditional expressions.

```
if      if_bit
elseif elseif_bit
while  while_bit
until  until_bit
for
```

[Generated instructions]

- (1) When the control statement is entered in lower case letters and CY has been entered

```
BNC    $?LFALSE
```

- (2) When the control statement is entered in lower case letters and Z has been entered

```
BNZ    $?LFALSE
```

- (3) When the control statement is entered in lower case letters and a bit symbol has been entered

```
BF     Bit symbol , $?LFALSE
```

- (4) When the control statement is entered in upper case letters and CY has been entered

```
      BC     $?LTRUE
      BR     ?LFALSE
?LTRUE :
```

- (5) When the control statement is entered in upper case letters and Z has been entered

```
      BZ     $?LTRUE
      BR     ?LFALSE
?LTRUE :
```

- (6) When the control statement is entered in upper case letters and a bit symbol has been entered.

BT	Bit symbol , \$?LTRUE
BR	?LFALSE
?LTRUE :	

[Use examples]

- (1) When the control statement is entered in lower case letters

Output source	Input source
BNC \$?L1	if_bit (CY)
CALL !XXX	CALL !XXX
BR ?L2	
?L1 :	else
CALL !YYY	CALL !YYY
?L2 :	endif
BNZ \$?L3	if_bit (Z)
CALL !XXX	CALL !XXX
BR ?L4	
?L3 :	else
CALL !YYY	CALL !YYY
?L4 :	endif
BF TRFG.0 , \$?L5	if_bit (TRFG.0)
CALL !XXX	CALL !XXX
BR ?L6	
?L5 :	else
CALL !YYY	CALL !YYY
?L6 :	endif

(2) When the control statement is entered in upper case letters

Output source	Input source
<pre> BC \$?L7 BR ?L8 ?L7 : </pre>	<pre> IF_BIT (CY) </pre>
<pre> CALL !XXX BR ?L9 ?L8 : </pre>	<pre> CALL !XXX ELSE CALL !YYY ENDIF </pre>
<pre> CALL !YYY ?L9 : </pre>	<pre> ENDIF </pre>
<pre> BZ \$?L10 BR ?L11 ?L10 : </pre>	<pre> IF_BIT (Z) </pre>
<pre> CALL !XXX BR ?L12 ?L11 : </pre>	<pre> CALL !XXX ELSE CALL !YYY ENDIF </pre>
<pre> CALL !YYY ?L12 : </pre>	<pre> ENDIF </pre>
<pre> BT TRFG.0 , \$?L13 BR ?L14 ?L13 : </pre>	<pre> IF_BIT (TRFG.0) </pre>
<pre> CALL !XXX BR ?L15 ?L14 : </pre>	<pre> CALL !XXX ELSE CALL !YYY ENDIF </pre>
<pre> CALL !YYY ?L15 : </pre>	<pre> ENDIF </pre>

Test bit expressions Negative logic (bit)

(2) !bit symbol

[Description format]

```
[ Δ ] !bit symbol [ Δ ]
```

[Function]

- True when the bit symbol contents are 0, false when they are 1.
- The following control statements are able to include bit symbols entered as conditional expressions.

```
if      if_bit
elseif elseif_bit
while  while_bit
until  until_bit
for
```

[Generated instructions]

- (1) When the control statement is entered in lower case letters and CY has been entered

```
BC      $?LFALSE
```

- (2) When the control statement is entered in lower case letters and Z has been entered

```
BZ      $?LFALSE
```

- (3) When the control statement is entered in lower case letters and a bit symbol has been entered

```
BT      Bit symbol , $?LFALSE
```

- (4) When the control statement is entered in upper case letters and CY has been entered

```
      BNC      $?LTRUE
      BR       ?LFALSE
?LTRUE :
```

- (5) When the control statement is entered in upper case letters and Z has been entered

```
      BNZ     $?LTRUE
      BR      ?LFALSE
?LTRUE :
```

- (6) When the control statement is entered in upper case letters and a bit symbol has been entered.

BF	Bit symbol , \$?LTRUE
BR	?LFALSE
?LTRUE :	

[Use examples]

- (1) When the control statement is entered in lower case letters

Output source	Input source
BC \$?L1	if_bit (!CY)
CALL !XXX	CALL !XXX
BR ?L2	
?L1 :	else
CALL !YYY	CALL !YYY
?L2 :	endif
	if_bit (!Z)
BZ \$?L3	CALL !XXX
CALL !XXX	
BR ?L4	
?L3 :	else
CALL !YYY	CALL !YYY
?L4 :	endif
	if_bit (!TRFG.0)
BT TRFG.0 , \$?L5	CALL !XXX
CALL !XXX	
BR ?L6	
?L5 :	else
CALL !YYY	CALL !YYY
?L6 :	endif

(2) When the control statement is entered in upper case letters

Output source	Input source
<pre> BNC \$?L7 BR ?L8 ?L7 : </pre>	<pre> IF_BIT (!CY) </pre>
<pre> CALL !XXX BR ?L9 ?L8 : </pre>	<pre> CALL !XXX ELSE CALL !YYY ENDIF </pre>
<pre> CALL !YYY ?L9 : </pre>	<pre> ENDIF </pre>
<pre> BNZ \$?L10 BR ?L11 ?L10 : </pre>	<pre> IF_BIT (!Z) </pre>
<pre> CALL !XXX BR ?L12 ?L11 : </pre>	<pre> CALL !XXX ELSE </pre>
<pre> CALL !YYY ?L12 : </pre>	<pre> CALL !YYY ENDIF </pre>
<pre> BF TRFG.0 , \$?L13 BR ?L14 ?L13 : </pre>	<pre> IF_BIT (!TRFG.0) </pre>
<pre> CALL !XXX BR ?L15 ?L14 : </pre>	<pre> CALL !XXX ELSE </pre>
<pre> CALL !YYY ?L15 : </pre>	<pre> CALL !YYY ENDIF </pre>

3.5.3 Logical operations

In the description of each type of conditional expression, it is noted that "?LTRUE" is used as the branch destination label when the test result is true and "?LFALSE" is used as this label when the test result is false.

A logical AND (&&) or logical OR (||) result can be obtained when there are two comparison expressions or a true/false test bit expression.

Up to 16 logical operators can be entered in a conditional expression.

This means that it is possible to enter expressions for processing that is executed when two conditional expressions are both met or when either of them are met.

The ST78K0S generates branch instructions beginning from the highest-priority logical operator.

(1) Code example

```
B < #0FFH && C >= #0 || D == #10
```

The following pages describe the functions of the various logical operations.

The use examples show as comment statements the source files to which generated instructions are input.

Table 3-6 Logical Operations

Logical Operation	Description Format	Function
Logical AND (&&)	Conditional expression 1 && conditional expression 2	True if both conditional expression 1 and conditional expression 2 are true
Logical OR ()	Conditional expression 1 conditional expression 2	True if either conditional expression 1 or conditional expression 2 is true

Logical operations Logical AND (&&)

(1) Logical AND (&&)

[Description format]

```
Conditional expression 1 [ Δ ] && [ Δ ] Conditional expression 2
```

[Function]

- The logical AND result of conditional expression 1 and conditional expression 2 is obtained. The result is true when conditional expression 1 and conditional expression 2 are both true and the result is false otherwise. The entered operation is performed when two conditions are met.

The output instruction differs depending on whether the control statement is entered in lower case letters or upper case letters.

Instructions for testing are generated first for contents enclosed in parentheses "(").

[Generated instructions]

- (1) When the control statement is entered in lower case letters

Table 3-7 Generated Instructions (Control Statement in Lower Case Letters) for Logical AND

Conditional expression	Generated instruction
$\alpha == \beta$ &&	CMP (W) α , β BNZ \$?LFALSE
$\alpha != \beta$ &&	CMP (W) α , β BZ \$?LFALSE
$\alpha < \beta$ &&	CMP (W) α , β BNC \$?LFALSE
$\alpha > \beta$ &&	CMP (W) α , β BZ \$?LFALSE BC \$?LFALSE
$\alpha >= \beta$ &&	CMP (W) α , β BC \$?LFALSE
$\alpha <= \beta$ &&	CMP (W) α , β BZ \$\$ + 4 BNZ \$?LFALSE
Bit symbol &&	BF Bit symbol , \$?LFALSE
CY &&	BNC \$?LFALSE
Z &&	BNZ \$?LFALSE
!bit symbol &&	BT Bit symbol , \$?LFALSE
!CY &&	BC \$?LFALSE
!Z &&	BZ \$?LFALSE

(2) When the control statement is entered in upper case letters

Table 3-8 Generated Instructions (Control Statement in Upper Case Letters) for Logical AND

Conditional expression	Generated instruction
$\alpha == \beta \ \&\&$	CMP (W) α , β BZ \$?LTRUE BR ?LFALSE ?LTRUE :
$\alpha != \beta \ \&\&$	CMP (W) α , β BNZ \$?LTRUE BR ?LFALSE ?LTRUE :
$\alpha < \beta \ \&\&$	CMP (W) α , β BC \$?LTRUE BR ?LFALSE ?LTRUE :
$\alpha > \beta \ \&\&$	CMP (W) α , β BZ \$\$ + 4 BNC \$?LTRUE BR ?LFALSE ?LTRUE :
$\alpha >= \beta \ \&\&$	CMP (W) α , β BNC \$?LTRUE BR ?LFALSE ?LTRUE :
$\alpha <= \beta \ \&\&$	CMP (W) α , β BZ \$?LTRUE BC \$?LTRUE BR ?LFALSE ?LTRUE :
Bit symbol $\&\&$	BT Bit symbol , \$?LTRUE BR ?LFALSE ?LTRUE :
CY $\&\&$	BC \$?LTRUE BR ?LFALSE ?LTRUE :
Z $\&\&$	BZ \$?LTRUE BR ?LFALSE ?LTRUE :
lbit symbol $\&\&$	BF Bit symbol , \$?LTRUE BR ?LFALSE ?LTRUE :
ICY $\&\&$	BNC \$?LTRUE BR ?LFALSE ?LTRUE :
IZ $\&\&$	BNZ \$?LTRUE BR ?LFALSE ?LTRUE :

[Use examples]

(1) When the control statement is entered in lower case letters

Output source	Input source
<pre> MOV A , C CMP A , #0 BNZ \$?L1 MOV A , B CMP A , #0 BC \$?L1 MOV A , B CMP A , #80H BNC \$?L1 CALL !XXX BR ?L2 ?L1 : CALL !YYY ?L2 :</pre>	<pre> if (C == #0 && B >= #0 && B < #80H) (A) CALL !XXX else CALL !YYY endif</pre>

(2) When the control statement is entered in upper case letters

Output source	Input source
<pre> MOV A , C CMP A , #0 BZ \$?L3 BR ?L6 ?L3 : MOV A , B CMP A , #0 BNC \$?L4 BR ?L6 ?L4 : MOV A , B CMP A , #80H BC \$?L5 BR ?L6 ?L5 : CALL !XXX BR ?L7 ?L6 : CALL !YYY ?L7 :</pre>	<pre> IF (C == #0 && B >= #0 && B < #80H) (A) CALL !XXX ELSE CALL !YYY ENDIF</pre>

Logical operations Logical OR (||)

(2) Logical OR (||)

[Description format]

```
Conditional expression 1 [ Δ ] || [ Δ ] Conditional expression 2
```

[Function]

- The logical OR result of conditional expression 1 and conditional expression 2 is obtained. The result is true when either conditional expression 1 or conditional expression 2 is true and the result is false when both are false. The entered operation is performed when either condition is met.

Instructions for testing are generated first for contents enclosed in parentheses "()".

[Generated instructions]

Table 3-9 Generated Instructions for Logical OR

Conditional expression	Generated instruction
$\alpha == \beta$	CMP (W) α , β BZ \$?LFALSE
$\alpha != \beta$	CMP (W) α , β BNZ \$?LFALSE
$\alpha < \beta$	CMP (W) α , β BC \$?LFALSE
$\alpha > \beta$	CMP (W) α , β BZ \$\$ + 4 BNC \$?LFALSE
$\alpha >= \beta$	CMP (W) α , β BNC \$?LFALSE
$\alpha <= \beta$	CMP (W) α , β BZ \$?LFALSE BC \$?LFALSE
Bit symbol	BT Bit symbol , \$?LFALSE
CY	BC \$?LFALSE
Z	BZ \$?LFALSE
!bit symbol	BF Bit symbol , \$?LFALSE
!CY	BNC \$?LFALSE
!Z	BNZ \$?LFALSE

[Use examples]

Output source	Input source
<pre> MOV A , B CMP A , #0 BZ \$?L1 MOV A , C CMP A , #0 BNC \$?L1 MOV A , D CMP A , #80H BNC \$?L2 ?L1 : CALL !XXX BR ?L3 ?L2 : CALL !YYY ?L3 :</pre>	<pre> if (B == #0 C >= #0 D < #80H) (A) CALL !XXX else CALL !YYY endif</pre>

CHAPTER 4 EXPRESSIONS

This chapter describes the functions of the expressions.

4.1 Overview

Expressions are used to perform assignments or arithmetic operations.

The following are examples of expressions

- Assignment statement : Assigns the second operand as the first operand
- Count statement : Adds or subtracts "1" to the operand value
- Exchange statement : Exchanges the values of the first and second operands
- Bit manipulation statement : Sets (to 1) or resets (to 0) the value of a operand

The functions of these expressions are described below.

The use examples show as comment statements the source files to which generated instructions are input.

Table 4-1 Assignment Statements

Assignment statement	Description format	Function
Assign (=)		
Assign	$\alpha = \beta$	$\alpha \leftarrow \beta$
Sequential assign	$\alpha_1 = \dots = \alpha_n = \beta$	$\alpha_1 \leftarrow \beta, \dots, \alpha_n \leftarrow \beta$
Assign (with register specification)	$\alpha = \beta (\gamma)$	$(\gamma) \leftarrow \beta, \alpha \leftarrow (\gamma)$
Sequential assign (with register specification)	$\alpha_1 = \dots = \alpha_n = \beta (\gamma)$	$\gamma \leftarrow \beta, \alpha_1 \leftarrow \gamma, \dots, \alpha_n \leftarrow \gamma$
IncrementAssign (+=)		
Increment assignment	$\alpha += \beta$	$\alpha \leftarrow \alpha + \beta$
Increment assignment (with register specification)	$\alpha += \beta (\text{register})$	$\gamma \leftarrow \alpha, \gamma \leftarrow \gamma + \beta, \alpha \leftarrow \gamma$
Increment assignment with carry	$\alpha += \beta , \text{CY}$	$\alpha \leftarrow \alpha + \beta, \text{CY}$
Increment assignment with carry (with register specification)	$\alpha += \beta , \text{CY} (\text{register})$	$\gamma \leftarrow \alpha, \gamma \leftarrow \gamma + \beta, \text{CY}, \alpha \leftarrow \gamma$

Table 4-1 Assignment Statements

Assignment statement	Description format	Function
DecrementAssign (-=)		
Decrement assignment	$\alpha -= \beta$	$\alpha \leftarrow \alpha - \beta$
Decrement assignment (with register specification)	$\alpha -= \beta , (\text{register})$	$\gamma \leftarrow \alpha, \gamma \leftarrow \gamma - \beta, \alpha \leftarrow \gamma$
Decrement assignment with carry	$\alpha -= \beta , \text{CY}$	$\alpha \leftarrow \alpha - \beta, \text{CY}$
Decrement assignment with carry (with register specification)	$\alpha -= \beta , \text{CY} (\text{register})$	$\gamma \leftarrow \alpha, \gamma \leftarrow \gamma - \beta, \text{CY}, \alpha \leftarrow \gamma$
LogicalANDAssign (&=)		
Logical AND assignment	$\alpha \&= \beta$	$\alpha \leftarrow \alpha \cap \beta$
Logical AND assignment (with register specification)	$\alpha \&= \beta (\text{register})$	$\gamma \leftarrow \alpha, \gamma \leftarrow \gamma \cap \beta, \alpha \leftarrow \gamma$
LogicalORAssign (=)		
Logical OR assignment	$\alpha = \beta$	$\alpha \leftarrow \alpha \cup \beta$
Logical OR assignment (with register specification)	$\alpha = \beta (\text{register})$	$\gamma \leftarrow \alpha, \gamma \leftarrow \gamma \cup \beta, \alpha \leftarrow \gamma$
LogicalXORAssign (^=)		
Logical XOR assignment	$\alpha ^= \beta$	$\alpha \leftarrow \alpha \wedge \beta$
Logical XOR assignment (with register specification)	$\alpha ^= \beta (\text{register})$	$\gamma \leftarrow \alpha, \gamma \leftarrow \gamma \wedge \beta, \alpha \leftarrow \gamma$
RightShiftAssign (>>=)		
Right shift (rotate) assignment	$\alpha \gg= \beta$	(α shifted to right of β bit)
Right shift assignment (with register specification)	$\alpha \gg= \beta (\text{register})$	$\gamma \leftarrow \alpha, (\gamma \text{ shifted to right of } \beta \text{ bit}), \alpha \leftarrow \gamma$
LeftShiftAssign (<<=)		
Left shift assignment	$\alpha \ll= \beta$	(α shifted to left of β bit)
Left shift assignment (with register specification)	$\alpha \ll= \beta (\text{register})$	$\gamma \leftarrow \alpha, (\gamma \text{ shifted to left of } \beta \text{ bit}), \alpha \leftarrow \gamma$

Table 4-2 Count Statements

Count statement	Description format	Function
Increment (++)	$\alpha ++$	$\alpha \leftarrow \alpha + 1$
Decrement (--)	$\alpha --$	$\alpha \leftarrow \alpha - 1$

Table 4-3 Exchange Statements

Exchange statement	Description format	Function
Exchange (<->)		
Exchange	$\alpha <-> \beta$	$\alpha <- \alpha <-> \beta$
Exchange (with register specification)	$\alpha <-> \beta (\gamma)$	$\gamma <- \alpha, \gamma <- \gamma <-> \beta, \alpha <- \gamma$

Table 4-4 Bit Manipulation Statements

Bit manipulation statement	Description format	Function
Set bit (=)		
Set bit	$\alpha = 1$	$\alpha <- 1$
Sequential set bit	$\alpha_1 = \dots = \alpha_n = 1$	$\alpha_n <- 1, \dots, \alpha_1 <- 1$
Set bit (with register specification)	$\alpha = 1 (CY)$	$CY <- 1, \alpha <- 1$
Sequential set bit (with register specification)	$\alpha_1 = \dots \alpha_n = 1 (CY)$	$CY <- 1, \alpha_n <- 1, \dots, \alpha_1 <- 1$
Clear bit (=)		
Clear bit	$\alpha = 0$	$\alpha <- 0$
Sequential clear bit	$\alpha_1 = \dots = \alpha_n = 0$	$\alpha_n <- 0, \dots, \alpha_1 <- 0$
Clear bit (with register specification)	$\alpha = 0 (CY)$	$CY <- 0, \alpha <- 0$
Sequential clear bit (with register specification)	$\alpha_1 = \dots \alpha_n = 0 (CY)$	$CY <- 0, \alpha_n <- 0, \dots, \alpha_1 <- 0$

4.2 Assignment Statements

Assignment statements Assign (=)

(1) Assign (=)

[Description format]

```
[ Δ ] [ size specification ] [ Δ ] α1 [ Δ ] [ = [ Δ ] [ size specification ] [ Δ ] α2 [ Δ ] ... ] = [ Δ ] [ size specification ] [ Δ ] β [ Δ ] [ ( register specification ) ]
```

[Function]

- When there is no register specification
β values on the right side are sequentially assigned to the left side.
- When there is a register specification
β values on the right side are assigned to the specified register or to CY and their contents are sequentially assigned to the left side.

[Description]

- α and β are values that can be entered via the MOV or MOVW instruction.
Up to 32 of the assignment operator "=" can be entered in one line. An error occurs when more than 32 are entered. If even one error occurs during sequential assignments, no instructions will be generated.

[Generated instructions](1) When α and β are not bit symbols<When β is CY>

	BF	β , ?L1
	SET1	CY
	BR	?L2
?L1 :	CLR1	CY
?L2 :		

However, sequential assignments cannot be entered.

<When β is CY>

	BNC	?L1
	SET1	α_n
	SET1	α_{n-1}
	:	
	SET1	α_2
	SET1	α_1
	BR	?L2
?L1 :	CLR1	α_n
	CLR1	α_{n-1}
	:	
	CLR1	α_2
	CLR1	α_1
?L2 :		

<When CY has been specified as the register>

	BF	β , ?L1
	SET1	α_n
	SET1	α_{n-1}
	:	
	SET1	α_2
	SET1	α_1
	BR	?L2
?L1 :	CLR1	α_n
	CLR1	α_{n-1}
	:	
	CLR1	α_2
	CLR1	α_1
?L2 :		

(2) When α and β are not bit symbols

<When there is no register specification>

```
MOV     $\alpha_1$  ,  $\beta$ 
```

MOV1 or MOVW may be generated instead, depending on the operand.

<When there is no register specification and a sequential assignment is entered>

```
MOV     $\alpha_n$  ,  $\beta$ 
MOV     $\alpha_{n-1}$  ,  $\beta$ 
      :
MOV     $\alpha_2$  ,  $\beta$ 
MOV     $\alpha_1$  ,  $\beta$ 
```

MOV1 or MOVW may be generated instead, depending on the operand.

<When there is a register specification>

```
MOV    Specified register ,  $\beta$ 
MOV     $\alpha_1$  , Specified register
```

MOV1 or MOVW may be generated instead, depending on the operand.

<When there is a register specification and a sequential assignment is entered>

```
MOV    Specified register ,  $\beta$ 
MOV     $\alpha_n$  , specified register
MOV     $\alpha_{n-1}$  , specified register
      :
MOV     $\alpha_2$  , specified register
MOV     $\alpha_1$  , specified register
```

MOV1 or MOVW may be generated instead, depending on the operand.

For details of combinations of α_n and β , see [Table 4-5](#). Depending on the entered statement, α_n and β , indicates the specified register.

[Use examples]

(1) When there is no register specification

Output source	Input source
<pre> BF P1.1 , \$?L1 SET1 CY BR ?L2 ?L1 : CLR1 CY ?L2 : MOV A , #4H MOVW AX , SYMP BNC \$?L3 SET1 bit1 SET1 PORT.0 BR ?L4 ?L3 : CLR1 bit1 CLR1 PORT.0 ?L4 : MOV DAT3 , A MOV DAT2 , A MOV DAT1 , A </pre>	<pre> CY = P1.1 A = #4H AX = SYMP PORT.0 = bit1 = CY DAT1 = DAT2 = DAT3 = A </pre>

(2) When there is a register specification

Output source	Input source
<pre> BF P1.1 , \$?L5 SET1 P0.2 SET1 A.0 BR ?L6 ?L5 : CLR1 P0.2 CLR1 A.0 ?L6 : MOV A , #4H MOV [DE] , A MOV A,X MOV DAT3 , A MOV DAT2 , A MOV DAT1 , A MOVW AX,BC MOVW DATA3P , AX MOVW DATA2P , AX MOVW DATA1P , AX </pre>	<pre> A.0 = P0.2 = P1.1 (CY) [DE] = #4H (A) DAT1 = DAT2 = DAT3 = X (A) DATA1P = DATA2P = DATA3P = BC (AX) </pre>

Table 4-5 Generated Instructions for Assignments

Symbol		β																					
		a	b	c	d	e	f	g	h	i	j	k	l	m	n	o	p	q	r	s	t	u	v
α_n	a	CY		*3		*4								*3									
	b	Bit symbol	*3			*5																	
	c	Byte user symbol	*3			*5																	
	d	[HL]. β	*6			*5		*1						*2									*1
	e	Byte data						*1															*1
	f	A				*1	*1		*1	*1										*1	*1	*1	*1
	g	Byte register						*1															*1
	h	R0						*1															*1
	i	R1																					*1
	j	sfr						*1															*1
	k	PSW						*1															*1
	l	Word user symbol	*3			*5		*1							*2								
	m	Word data													*2								
	n	AX				*2								*2		*2		*2		*2			*2
	o	Word register													*2								*2
	p	RP0																					*2
	q	sfrp																					
	r	SP													*2								
	s	Direct access symbol						*1															
	t	Indirect access symbol						*1															
	u	[DE]						*1															
	v	Immediate symbol																					

*1 : Generates MOV instruction

*2 : Generates MOVW instruction

*3 : Generates MOV1 instruction

*4 : Generates SET1 instruction when "1" has been entered as b. Generates CLR1 instruction when "0" has been entered. Generates MOV when any value other than "0" or "1" has been entered.

*5 : Generates SET1 when "1" has been entered as b. Generates CLR1 when "0" has been entered.

*6 : Generates MOV1 when any value other than "0" or "1" has been entered as an

Empty spaces indicate errors.

Assignment statements IncrementAssign (+=)

(2) IncrementAssign (+=)

[Description format]

```
[ Δ ] [ size specification ] [ Δ ] α [ Δ ] += [ Δ ] [ size specification ] [ Δ ] β [ Δ ] [ , [ Δ ] CY ] [ Δ ] [ ( register specification ) ]
```

[Function]

- When there is no register specification
The two operands α and β are added and the result is assigned to α .
- When there is a register specification
 α is assigned to the specified register.
The contents of the specified register are added to β and their result is assigned to the specified register.
The contents of the specified register are assigned to α .
- Increment with carry ; no register specification
An increment with carry operation is performed using the two operands α and β , and the result is assigned to α .
- Increment with carry ; with register specification
The contents of α are assigned to the specified register.
An increment with carry operation is performed using the contents of the specified register and β , and the result is assigned to the specified register.
The contents of the specified register are assigned to α .

[Description]

- When there is no register specification
The contents of α and β can be entered in ADD and ADDW.
- When there is a register specification
The contents of α can be entered in MOV and MOVW.
The contents of β can be entered in ADD and ADDW.
- Increment with carry ; no register specification
The contents of α and β can be entered in ADDC.
- Increment with carry ; with register specification
The contents of α can be entered in MOV.
The contents of β can be entered in ADDC.

[Generated instructions]

- (1) When there is no register specification

ADD	α , β
-----	--------------------

ADDW may be generated instead, depending on the operand.

- (2) When there is a register specification

MOV	Specified register , α
ADD	Specified register , β
MOV	α , specified register

ADDW may be generated instead, depending on the operand.

- (3) Increment with carry ; no register specification

ADDC	α , β
------	--------------------

- (4) Increment with carry ; with register specification

MOV	Specified register , α
ADDC	Specified register , β
MOV	α , specified register

For details of combinations of α and β , see [Table 4-6](#). Depending on the entered statement, α indicates the specified register.

[Use examples]

- (1) When there is no register specification

Output source	Input source
ADD A , #0C0H ADDW AX , #0C00H	A += #0C0H AX += #0C00H

- (2) When there is a register specification

Output source	Input source
MOV A , !ABC ADD A , #0FCH MOV !ABC , A MOVW AX , HL ADDW AX , #0FFFH MOVW HL , AX	ABC += #0FCH (A) HL += #0FFFH (AX)

(3) Increment with carry ; no register specification

Output source	Input source
ADDC A , #50H	A += #50H , CY

(4) Increment with carry ; with register specification

Output source	Input source
MOV A , PSW ADDC A , #50H MOV PSW , A	PSW += #50H , CY (A)

Table 4-6 Generated Instructions for Increment Assignments

シンボル	β																										
	a	b	c	d	e	f	g	h	i	j	k	l	m	n	o	p	q	r	s	t	u	v					
α	a	CY																									
	b	Bit symbol																									
	c	Byte user symbol																									
	d	[HL].β																								*1	
	e	Byte data																								*1	
	f	A			*1	*1	*1	*1	*1			*1							*1	*1						*1	
	g	Byte register																									
	h	R0																									
	i	R1																									
	j	sfr																									
	k	PSW																									
	l	Word user symbol																									
	m	Word data																									
	n	AX																								*2	
	o	Word register																									
	p	RP0																									
	q	sfrp																									
	r	SP																									
	s	Direct access symbol																									
	t	Indirect access symbol																									
	u	[DE]																									
	v	Immediate symbol																									

*1 : Generates ADD instruction. For increment with carry, ADDC instruction is generated.

*2 : Generates ADDW instruction.

Empty spaces indicate errors.

Assignment statements DecrementAssign (-=)

(3) DecrementAssign (-=)

[Description format]

```
[ Δ ] [ size specification ] [ Δ ] α [ Δ ] -= [ Δ ] [ size specification ] [
Δ ] β [ Δ ] [ , [ Δ ] CY ] [ Δ ] [ ( register specification ) ]
```

[Function]

- When there is no register specification
β is subtracted from α and the result is assigned to α.
- When there is a register specification
α is assigned to the specified register.
β is subtracted from the contents of the specified register and the result is assigned to the specified register.
The contents of the specified register are assigned to α.
- Decrement with carry ; no register specification
A decrement with carry operation is performed using the two operands α and β, and the result is assigned to α.
- Decrement with carry ; with register specification
The contents of α are assigned to the specified register.
An decrement with carry operation is performed using the contents of the specified register and β, and the result is assigned to the specified register.
The contents of the specified register are assigned to α.

[Description]

- When there is no register specification
The contents of α and β can be entered in SUB and SUBW.
- When there is a register specification
The contents of α can be entered in MOV and MOVW.
The contents of β can be entered in SUB and SUBW.
- Decrement with carry ; no register specification
The contents of α and β can be entered in SUBC.
- Decrement with carry ; with register specification
The contents of α can be entered in MOV.
The contents of β can be entered in SUBC.

[Generated instructions]

- (1) When there is no register specification

```
SUB    α , β
```

SUBW may be generated instead, depending on the operand.

- (2) When there is a register specification

```
MOV    Specified register , α
SUB    Specified register , β
MOV    α , specified register
```

SUBW may be generated instead, depending on the operand.

- (3) Decrement with carry ; no register specification

```
SUBC   α , β
```

- (4) Decrement with carry ; with register specification

```
MOV    Specified register , α
SUBC   Specified register , β
MOV    α , specified register
```

For details of combinations of α and β , see [Table 4-7](#). Depending on the entered statement, α indicates the specified register.

[Use examples]

- (1) When there is no register specification

Output source	Input source
SUB A , #0C0H SUBW AX , #0C00H	A -= #0C0H AX -= #0C00H

- (2) When there is a register specification

Output source	Input source
MOV A , !ABC SUB A , #0FCH MOV !ABC , A MOVW AX , HL SUBW AX , #0FFFH MOVW HL , AX	!ABC -= #0FCH (A) HL -= #0FFFH (AX)

(3) Decrement with carry ; no register specification

Output source	Input source
SUBC A , #50H	A -= #50H , CY

(4) Decrement with carry ; with register specification

Output source	Input source
MOV A , PSW SUBC A , #50H MOV PSW , A	PSW -= #50H , CY (A)

Table 4-7 Generated Instructions for Decrement Assignments

Symbol	β																										
	a	b	c	d	e	f	g	h	i	j	k	l	m	n	o	p	q	r	s	t	u	v					
α a	CY																										
b	Bit symbol																										
c	Byte user symbol																										
d	[HL]. β																									*1	
e	Byte data																										*1
f	A				*1	*1		*1	*1	*1			*1							*1	*1						*1
g	Byte register																										
h	R0																										
i	R1																										
j	sfr																										
k	PSW																										
l	Word user symbol																										
m	Word data																										
n	AX																										*2
o	Word register																										
p	RP0																										
q	sfrp																										
r	SP																										
s	Direct access symbol																										
t	Indirect access symbol																										
u	[DE]																										
v	Immediate symbol																										

*1 : Generates SUB instruction. For decrement with carry, SUBC instruction is generated.

*2 : Generates SUBW instruction.

Empty spaces indicate errors.

Assignment statements LogicalANDAssign (&=)

(4) LogicalANDAssign (&=)

[Description format]

```
[ Δ ] [ size specification ] [ Δ ] α [ Δ ] &= [ Δ ] [ size specification ] [ Δ ] β [ Δ ] [ register specification ]
```

[Function]

- When there is no register specification

The logical AND (α & β) is obtained from the bits in α and β , and the result is assigned to α .

- When there is a register specification

α is assigned to the specified register.

The logical AND (specified register & β) is obtained from the bits in the specified register and β , and the result is assigned to the specified register.

The contents of the specified register are assigned to α .

[Description]

- Where there is no register specification

The contents of α and β can be entered in AND and BF.

- Where there is a register specification

The contents of α can be entered in MOV and BF.

The contents of β can be entered in AND and BF.

[Generated instructions]

- (1) When there is no register specification

< When α is CY >

```
BNC      ?L1
BF       β , ?L1
SET1     CY
BR       ?L2
```

?L1 :

```
CLR1     CY
```

?L2 :

< When α is not CY >

```
AND      α , β
```

(2) When there is a register specification

< When the specified register is CY >

	BF	α , ?L1
	BF	β , ?L1
	SET1	α
	BR	?L2
?L1 :	CLR1	α
?L2 :		

< When the specified register is not CY >

MOV	Specified register	, α
AND	Specified register	, β
MOV	α	, specified register

For details of combinations of α and β , see [Table 4-8](#).

[Use examples]

(1) When there is no register specification

Output source	Input source
BNC $\$?L1$	CY &= P1S.1
BF P1S.1 , $\$?L1$	
SET1 CY	
BR ?L2	
?L1 : CLR1 CY	A &= #0FFH
?L2 : AND A , #0FFH	

(2) When there is a register specification

Output source	Input source
BF A.1 , $\$?L3$	A.1 &= PORT3.0 (CY)
BF PORT3.0 , $\$?L3$	
SET1 A.1	
BR ?L4	
?L3 : CLR1 A.1	[DE] &= #07H (A)
?L4 : MOV A , [DE]	
AND A , #07H	
MOV [DE] , A	

Table 4-8 Generated Instructions for Logical AND Assignments

シンボル		β																							
		a	b	c	d	e	f	g	h	i	j	k	l	m	n	o	p	q	r	s	t	u	v		
α	a	CY	*2	*2								*2													
	b	Bit symbol																							
	c	Byte user symbol																							
	d	[HL].β																						*1	
	e	Byte data																							*1
	f	A			*1	*1		*1	*1	*1			*1							*1	*1				*1
	g	Byte register																							
	h	R0																							
	i	R1																							
	j	sfr																							
	k	PSW																							
	l	Word user symbol																							
	m	Word data																							
	n	AX																							
	o	Word register																							
	p	RP0																							
	q	sfrp																							
	r	SP																							
	s	Direct access symbol																							
	t	Indirect access symbol																							
u	[DE]																								
v	Immediate symbol																								

*1 : Generates AND instruction.

*2 : Generates AND1 instruction.

Empty spaces indicate errors.

Assignment statements LogicalORAssign (|=)

(5) LogicalORAssign (|=)

[Description format]

```
[ Δ ] [ size specification ] [ Δ ] α [ Δ ] |= [ Δ ] [ size specification ] [ Δ ] β [ Δ ] [ register specification ]
```

[Function]

- When there is no register specification

The logical OR (α | β) is obtained from the bits in α and β , and the result is assigned to α .

- When there is a register specification

α is assigned to the specified register.

The logical OR (specified register | β) is obtained from the bits in the specified register and β , and the result is assigned to the specified register.

The contents of the specified register are assigned to α .

[Description]

- When there is no register specification

The contents of α and β can be entered in OR and BF.

- When there is a register specification

The contents of α can be entered in MOV and BT.

The contents of β can be entered in OR and BF.

[Generated instructions]

- (1) When there is no register specification

< When α is CY >

	BC	?L1
	BF	β , ?L2
?L1 :	SET1	CY
	BR	?L3
?L2 :	CLR1	CY
?L3 :		

< When α is not CY >

OR	α , β
----	--------------------

(2) When there is a register specification

< When the specified register is CY >

	BC	?L1
	BF	β , ?L2
?L1 :	SET1	CY
	BR	?L3
?L2 :	CLR1	CY
?L3 :		

< When the specified register is not CY >

MOV	Specified register , α
OR	Specified register , β
MOV	α , specified register

For details of combinations of α and β , see [Table 4-9](#).**[Use examples]**

(1) When there is no register specification

Output source		Input source
	BC $\$?L1$	CY = P1S.1
	BF P1S.1 , $\$?L2$	
?L1 :	SET1 CY	A = #0FFH
	BR ?L3	
?L2 :	CLR1 CY	
?L3 :	OR A , #0FFH	

(2) When there is a register specification

Output source		Input source
	BT A.1 , $\$?L4$	A.1 = PORT3.0 (CY)
	BF PORT3.0 , $\$?L5$	
?L4 :	SET1 A.1	[DE] = #07H (A)
	BR ?L6	
?L5 :	CLR1 A.1	
?L6 :	MOV A , [DE]	
	OR A , #07H	
	MOV [DE] , A	

Table 4-9 Generated Instructions for Logical OR Assignments

Symbol		β																						
		a	b	c	d	e	f	g	h	i	j	k	l	m	n	o	p	q	r	s	t	u	v	
α	a	CY	*2	*2								*2												
	b	Bit symbol																						
	c	Byte user symbol																						
	d	[HL]. β																					*1	
	e	Byte data																						*1
	f	A			*1	*1		*1	*1	*1			*1							*1	*1			*1
	g	Byte register																						
	h	R0																						
	i	R1																						
	j	sfr																						
	k	PSW																						
	l	Word user symbol																						
	m	Word data																						
	n	AX																						
	o	Word register																						
	p	RP0																						
	q	sfrp																						
	r	SP																						
	s	Direct access symbol																						
	t	Indirect access symbol																						
	u	[DE]																						
	v	Immediate symbol																						

*1 : Generates OR instruction.

*2 : Generates OR1 instruction.

Empty spaces indicate errors.

Assignment statements LogicalXORAssign (^=)

(6) LogicalXORAssign (^=)

[Description format]

```
[ Δ ] [ size specification ] [ Δ ] α [ Δ ] ^= [ Δ ] [ size specification ] [ Δ ] β [ Δ ] [ register specification ]
```

[Function]

- When there is no register specification
The logical XOR ($\alpha \wedge \beta$) is obtained from the bits in α and β , and the result is assigned to α .
- When there is a register specification
 α is assigned to the specified register.
The logical XOR (specified register $\wedge \beta$) is obtained from the bits in the specified register and β , and the result is assigned to the specified register.
The contents of the specified register are assigned to α .

[Description]

- When there is no register specification
The contents of α and β can be entered in XOR and BF.
- When there is a register specification
The contents of α can be entered in MOV and BT.
The contents of β can be entered in XOR and BF.

[Generated instructions]

- (1) When there is no register specification

< When α is CY >

	BNC	?L1
	BF	β , ?L2
?L1 :	BC	?L3
	BF	β , ?L3
?L2 :	SET1	CY
	BR	?L4
?L3 :	CLR1	CY
?L4 :		

< When α is not CY >

XOR	α , β
-----	--------------------

(2) When there is a register specification

< When the specified register is CY >

	BF	α , ?L1
	BF	β , ?L2
?L1 :		
	BT	α , ?L3
	BF	β , ?L3
?L2 :		
	SET1	α
	BR	?L4
?L3 :		
	CLR1	α
?L4 :		

< When the specified register is not CY >

MOV	Specified register , α
XOR	Specified register , β
MOV	α , specified register

For details of combinations of α and β , see [Table 4-10](#).

[Use examples]

(1) When there is no register specification

Output source		Input source
	BNC \$?L1	CY ^= P1S.1
	BF P1S.1 , \$?L2	
?L1 :		
	BC \$?L3	
	BF P1S.1 , \$?L3	
?L2 :		
	SET1 CY	
	BR ?L4	
?L3 :		
	CLR1 CY	
?L4 :		
	XOR A , #0FFH	A ^= #0FFH

(2) When there is a register specification

Output source			Input source	
	BF	A.1 , \$?L5	A.1 ^= PORT3.0 (CY)	
	BF	PORT3.0 , \$?L6		
?L5 :				
	BT	A.1 , \$?L7		
	BF	PORT3.0 , \$?L7		
?L6 :				
	SET1	A.1		
	BR	?L8		
?L7 :				
	CLR1	A.1		
?L8 :				
	MOV	A , [DE]		[DE] ^= #07H (A)
	XOR	A , #07H		
	MOV	[DE] , A		

Table 4-10 Generated Instructions for Logical XOR Assignments

Symbol		β																											
		a	b	c	d	e	f	g	h	i	j	k	l	m	n	o	p	q	r	s	t	u	v						
α	a	CY	*2	*2								*2																	
	b	Bit symbol																											
	c	Byte user symbol																											
	d	[HL]. β																										*1	
	e	Byte data																										*1	
	f	A				*1	*1		*1	*1	*1		*1									*1	*1					*1	
	g	Byte register																											
	h	R0																											
	i	R1																											
	j	sfr																											
	k	PSW																											
	l	Word user symbol																											
	m	Word data																											
	n	AX																											
	o	Word register																											
	p	RP0																											
	q	sfrp																											
	r	SP																											
	s	Direct access symbol																											
	t	Indirect access symbol																											
	u	[DE]																											
	v	Immediate symbol																											

*1 : Generates XOR instruction.

*2 : Generates XOR1 instruction.

Empty spaces indicate errors.

Assignment statements RightShiftAssign (>>=)

(7) RightShiftAssign (>>=)

[Description format]

```
[ Δ ] [ size specification ] [ Δ ] α [ Δ ] >>= [ Δ ] β [ Δ ] [ ( register
specification ) ]
```

[Function]

- When there is no register specification
α is shifted to the right of the β bit, and the result is assigned to α.
- When there is a register specification
α is assigned to the specified register.
The contents of the specified register are shifted to the right of the β bit, and the result is assigned to the specified register.
The contents of the specified register are assigned to α.

[Description]

- When there is no register specification
The contents of α can be entered in A only.
The contents of β can be entered as numerals from 1 to 7.
- When there is a register specification
The contents of α can be entered in MOV.
The contents of β can be entered as numerals from 1 to 7.
The specified register can be entered in A only.

[Generated instructions]

- (1) When there is no register specification
An AND instruction is generated after a ROR instruction is output β times.

```
ROR    A , 1
      :
AND    A , #0FFH SHR β
```

- (2) When there is a register specification

```
MOV    A , α
ROR    A , 1
      :
AND    A , #0FFH SHR β
MOV    α , A
```

[Use examples]

(1) When there is no register specification

Output source	Input source
ROR A , 1 ROR A , 1 ROR A , 1 ROR A , 1 AND A , #0FFH SHR 4	A >>= 4

(2) When there is a register specification

Output source	Input source
MOV A , CCV ROR A , 1 ROR A , 1 ROR A , 1 ROR A , 1 AND A , #0FFH SHR 4 MOV CCV , A	CCV >>= 4 (A)

Assignment statements LeftShiftAssign (<<=)

(8) LeftShiftAssign (<<=)

[Description format]

```
[ Δ ] [ size specification ] [ Δ ] α [ Δ ] <<= [ Δ ] β [ Δ ] [ ( register specification ) ]
```

[Function]

- When there is no register specification
α is shifted to the left of the β bit, and the result is assigned to α.
- When there is a register specification
α is assigned to the specified register.
The contents of the specified register are shifted to the left of the β bit, and the result is assigned to the specified register.
The contents of the specified register are assigned to α.

[Description]

- When there is no register specification
The contents of α can be entered in A only.
The contents of β can be entered as numerals from 1 to 7.
- When there is a register specification
The contents of α can be entered in MOV.
The contents of β can be entered as numerals from 1 to 7.
The specified register can be entered in A only.

[Generated instructions]

- (1) When there is no register specification
An AND instruction is generated after a ROL instruction is output β times.

```
ROL    A , 1
      :
AND    A , #LOW ( 0FFH SHL β )
```

- (2) When there is a register specification

```
MOV    A , α
ROL    A , 1
      :
AND    A , #LOW ( 0FFH SHL β )
MOV    α , A
```

[Use examples]

(1) When there is no register specification

Output source	Input source
ROL A , 1 ROL A , 1 ROL A , 1 ROL A , 1 AND A , #LOW (0FFH SHL 4)	A <<= 4

(2) When there is a register specification

Output source	Input source
MOV A , CCV ROL A , 1 ROL A , 1 ROL A , 1 ROL A , 1 AND A , #LOW (0FFH SHL 4) MOV CCV , A	CCV <<= 4 (A)

4.3 Count Statements

Count statements Increment (++)

(9) Increment (++)

[Description format]

```
[ Δ ] [ size specification ] [ Δ ] α [ Δ ] ++
```

[Function]

- 1 is added to the contents of α .

[Description]

- The contents of α can be entered in INC or INCW.

[Generated instructions]

```
INC    α
```

INCW may be generated depending on the operands.

For details of α , see [Table 4-11](#).

[Use examples]

Output source		Input source	
INC	H	H++	
INC	CNT	CNT++	
INCW	HL	HL++	

Table 4-11 Generated Instructions for Increment

Symbol			Generated instructions
α	a	CY	
	b	Bit symbol	
	c	[HL]. β	
	d	Byte user symbol	*1
	e	Byte data	*1
	f	A	*1
	g	Byte register	*1
	h	R0	*1
	i	R1	*1
	j	sfr	
	k	PSW	
	l	Word user symbol	
	m	Word data	
	n	AX	*2
	o	Word register	*2
	p	RP0	*2
	q	sfrp	
	r	SP	
	s	Direct access symbol	
	t	Indirect access symbol	
	u	[DE]	
	v	Immediate symbol	

*1 : Generates INC instruction.

*2 : Generates INCW instruction.

Empty spaces indicate errors.

Count statements Decrement (--)

(10) Decrement (--)

[Description format]

```
[ Δ ] [ size specification ] [ Δ ] α [ Δ ] --
```

[Function]

- 1 is subtracted from the contents of α .

[Description]

- The contents of α can be entered in DEC or DECW.

[Generated instructions]

```
DEC    α
```

DECW may be generated depending on the operands.

For details of a, see [Table 4-12](#).

[Use examples]

Output source	Input source
DEC H	H--
DEC CNT	CNT--
DECW HL	HL--

Table 4-12 Generated Instructions for Decrement

Symbol			Generated instructions
α	a	CY	
	b	Bit symbol	
	c	[HL]. β	
	d	Byte user symbol	*1
	e	Byte data	*1
	f	A	*1
	g	Byte register	*1
	h	R0	*1
	i	R1	*1
	j	sfr	
	k	PSW	
	l	Word user symbol	
	m	Word data	
	n	AX	*2
	o	Word register	*2
	p	RP0	*2
	q	sfrp	
	r	SP	
	s	Direct access symbol	
	t	Indirect access symbol	
u	[DE]		
v	Immediate symbol		

*1 : Generates DEC instruction.

*2 : Generates DECW instruction.

Empty spaces indicate errors.

4.4 Exchange Statements

Exchange statements Exchange (<->)

(11) Exchange (<->)

[Description format]

```
[ Δ ] [ size specification ] [ Δ ] α [ Δ ] <-> [ Δ ] [ size specification ] [ Δ ] β [ Δ ] [ ( register specification ) ]
```

[Function]

- When there is no register specification
The contents of α and β are exchanged.
- When there is a register specification
The contents of α are assigned to the specified register.
The contents of the specified register are exchanged with the contents of β .
The contents of the specified register are assigned to α .

[Description]

- Where there is no register specification
The contents of α and β can be entered in XCH or XCHW.
- When there is a register specification
The contents of α can be entered in MOV and MOVW.
The contents of β can be entered in XCH and XCHW.

[Generated instructions]

- (1) When there is no register specification

```
XCH    α , β
```

XCHW may be generated depending on the operands.

- (2) When there is a register specification

```
MOV    Specified register , α
XCH    Specified register , β
MOV    α , specified register
```

XCHW may be generated depending on the operands.

For details of combinations of α and β , see [Table 4-13](#).

α indicates the specified register.

[Use examples]

- (1) When there is no register specification

Output source	Input source
XCH A , B XCHW AX , BC	A <-> B AX <-> BC

- (2) When there is a register specification

Output source	Input source
MOV A , DATA XCH A , B MOV DATA , A MOVW AX , DE XCHW AX , BC MOVW DE , AX	DATA <-> B (A) DE <-> BC (AX)

Table 4-13 Generated Instructions for Exchange

Symbol		β																					
		a	b	c	d	e	f	g	h	i	j	k	l	m	n	o	p	q	r	s	t	u	v
α	a	CY																					
	b	Bit symbol																					
	c	Byte user symbol																					
	d	[HL]. β																					
	e	Byte data																					
	f	A			*1	*1		*1			*1		*1									*1	*1
	g	Byte register																					
	h	R0																					
	i	R1																					
	j	sfr																					
	k	PSW																					
	l	Word user symbol																					
	m	Word data																					
	n	AX																					
	o	Word register																					
	p	RP0																					
	q	sfrp																					
	r	SP																					
	s	Direct access symbol																					
	t	Indirect access symbol																					
	u	[DE]																					
	v	Immediate symbol																					

*1 : Generates XCH instructions.

*2 : Generates XCHW instructions.

Empty spaces indicate errors.

4.5 Bit Manipulation Statements

Bit manipulation statements Set bit (=)

(12) Set bit (=)

[Description format]

```
[ Δ ] α1 [ Δ ] [ = [ Δ ] α2 [ Δ ] ... ] = [ Δ ] 1 [ Δ ] [ ( CY specification ) ]
Enter a "1" at the end of the right side.
```

[Function]

- When there is no CY specification
α_n is set (to a value of "1").
- When there is a CY specification
CY and α_n are set (to a value of "1").

[Description]

- The contents of α_n can be entered in a SET1 instruction.
Up to 32 of the assignment operator "=" can be entered in one line. An error occurs when more than 32 are entered. If even one error occurs during sequential assignments, no instructions will be generated.

[Generated instructions]

- (1) When there is no CY specification

```
SET1    α1
```

- (2) When there is no CY specification in sequential assignments

```
SET1    αn
SET1    αn-1
        :
SET1    α2
SET1    α1
```

- (3) When there is a CY specification

```
SET1    CY
SET1    α1
```

- (4) When there is a CY specification in sequential assignments

SET1	CY
SET1	α_n
SET1	α_{n-1}
	:
SET1	α_2
SET1	α_1

For details, see [Table 4-14](#).

[Use examples]

- (1) When there is no CY specification

Output source	Input source
SET1 A.3	A.3 = 1
SET1 CY	CY = 1
SET1 BIT3	BIT1 = BIT2 = BIT3 = 1
SET1 BIT2	
SET1 BIT1	

- (2) When there is a CY specification

Output source	Input source
SET1 CY	A.5 = 1 (CY)
SET1 A.5	
SET1 CY	BIT1 = BIT2 = BIT3 = 1 (CY)
SET1 BIT3	
SET1 BIT2	
SET1 BIT1	

Table 4-14 Generated Instructions for Set Bit

Symbol			Generated instructions
α	a	CY	*1
	b	Bit symbol	*1
	c	[HL]. β	*1
	d	Byte user symbol	*1
	e	Byte data	
	f	A	
	g	Byte register	
	h	R0	
	i	R1	
	j	sfr	
	k	PSW	
	l	Word user symbol	*1
	m	Word data	
	n	AX	
	o	Word register	
	p	RP0	
	q	sfrp	
	r	SP	
	s	Direct access symbol	
	t	Indirect access symbol	
u	[DE]		
v	Immediate symbol		

*1 : Generates SET1 instruction.

Empty spaces indicate errors.

Bit manipulation statements Clear bit (=)

(13) Clear bit (=)

[Description format]

```
[ Δ ] α1 [ = [ Δ ] α2 [ Δ ] ... ] = [ Δ ] 0 [ Δ ] [ ( CY specification ) ]
Enter a "0" at the end of the right side.
```

[Function]

- When there is no CY specification
α_n is cleared (to a value of "0").
- When there is a CY specification
CY and α_n are cleared (to a value of "0").

[Description]

- The contents of α_n can be entered in a CLR1 instruction.
Up to 32 of the assignment operator "=" can be entered in one line. An error occurs when more than 32 are entered. If even one error occurs during sequential assignments, no instructions will be generated.

[Generated instructions]

- (1) When there is no CY specification

```
CLR1    α1
```

- (2) When there is no CY specification in sequential assignments

```
CLR1    αn
CLR1    αn-1
        :
CLR1    α2
CLR1    α1
```

- (3) When there is a CY specification

```
CLR1    CY
CLR1    α1
```

- (4) When there is a CY specification in sequential assignments

CLR1	CY
CLR1	α_n
CLR1	α_{n-1}
	:
CLR1	α_2
CLR1	α_1

For details, see [Table 4-15](#).

[Use examples]

- (1) When there is no CY specification

Output source	Input source
CLR1 A.3 CLR1 CY CLR1 BIT3 CLR1 BIT2 CLR1 BIT1	A.3 = 0 CY = 0 BIT1 = BIT2 = BIT3 = 0

- (2) When there is a CY specification

Output source	Input source
CLR1 CY CLR1 A.5 CLR1 CY CLR1 BIT3 CLR1 BIT2 CLR1 BIT1	A.5 = 0 (CY) BIT1 = BIT2 = BIT3 = 0 (CY)

Table 4-15 Generated Instructions for Clear Bit

Symbol			Generated instructions
α	a	CY	*1
	b	Bit symbol	*1
	c	[HL]. β	*1
	d	Byte user symbol	*1
	e	Byte data	
	f	A	
	g	Byte register	
	h	R0	
	i	R1	
	j	sfr	
	k	PSW	
	l	Word user symbol	*1
	m	Word data	
	n	AX	
	o	Word register	
	p	RP0	
	q	sfrp	
	r	SP	
	s	Direct access symbol	
	t	Indirect access symbol	
u	[DE]		
v	Immediate symbol		

*1 : Generates CLR1 instruction.

Empty spaces indicate errors.

CHAPTER 5 DIRECTIVES

This chapter describes directives. In this case, "directives" means various directives that the ST78K0S requires to execute a series of processes.

5.1 Overview

Directives are entered into source programs as various directives that the ST78K0S requires to execute a series of processes.

The use of directives can make source program coding easier.

Directives are not output in output files.

5.2 Directive Functions

The following pages describe the functions of the various directive functions.

The use examples show as comment statements the source files to which generated instructions are input.

Table 5-1 List of Directives

Type of directive	Directive name
Symbol definition directive (#define)	#define
Conditional processing directive (#ifdef / #else / #endif)	#ifdef : #else : #endif
Include directive (#include)	#include
CALLT replacement directive (#defcallt)	#defcallt : #endcallt

#DEFINE

(1) Symbol definition directive (#define)

[Description format]

```
[ Δ ] # [ Δ ] define Δ symbol Δ character string
```

[Function]

- This directive replaces the specified character string with a symbol that has been entered in the source program.

[Description]

- The "#" character must always be entered at the start of the symbol, except when starting with a white space or a horizontal tab.
- Symbols start with a letter and are composed of alphanumeric characters. The first 31 characters are valid. If a symbol with more than 31 characters is specified, the 32nd and subsequent characters are ignored.
- Character strings are defined as strings of characters from among the characters in the set listed in ["2.2 \(1\) Character set"](#). They cannot include white spaces or quotation marks. Any character strings that contain white spaces or quotation marks will be ignored as processing continues.
- This directive is useful when coding easy-to-read symbols, such as numerical values.
- Reserved words cannot be entered as symbols.
- Reserved words can be entered as character strings.
- If the same symbol is defined twice, a warning message is output.
- Character strings that have been converted to secondary source files are output. The #define statement is not output.
- If a converted character string has already been defined by another #define statement, it can be reconverted up to 31 times. An error message is output during the 32nd conversion, and the definition is ignored during subsequent conversions.
- This directive can be entered anywhere in the source code.
- A warning message is output when two or more symbols specifying option "-d" are entered, and the #define statement is valid.

[Use examples]

Output source	Input source
<pre> MOV X , #0 CALL !xxx MOV A , X CMP A , #TRUE BNZ \$?L1 MOV B , #0C5H ?L1 :</pre>	<pre> #define TRUE 1 X = #0 CALL !xxx if (X == #TRUE) (A) B = #0C5H endif</pre>

#IFDEF / #ELSE / #ENDIF

(2) Conditional processing directive (#ifdef / #else / #endif)

[Description format]

```
[ Δ ] # [ Δ ] ifdef Δ symbol
      text 1
[ Δ ] # [ Δ ] else
      text 2
[ Δ ] # [ Δ ] endif
```

[Function]

- This directive performs conditional processing.
 - (1) When the symbol has not been defined
 - If #else has been entered, text 1 is skipped and text 2 becomes a processing object.
 - (2) When the symbol has been defined
 - If #else has been entered, text 1 becomes a processing object and text 2 is skipped.

[Description]

- The "#" character must always be entered at the start of the symbol, except when starting with a white space or a horizontal tab.
- Symbols start with a letter and are composed of alphanumeric characters. The first 31 characters are valid.
- Symbols are defined by a previously entered #define statement or by specifying the "-d" option at startup.
- This directive can be nested in up to eight levels.
- #else can be omitted.

[Use examples]

- When the following has been entered on the command line (and the symbol has been defined)

C > st78k0s -cP9014 sample.st -dSYM

Output source	Input source
MOV A , #00H	#ifdef SYM A = #00H #else A = #0FFH #endif

- When the following has been entered on the command line (and the symbol has not been defined)

C > st780 -cP9014 sample.st

Output source	Input source
MOV A , #0FFH	#ifdef SYM A = #00H #else A = #0FFH #endif

#INCLUDE

(3) Include directive (#include)

[Description format]

```
[ Δ ] # [ Δ ] include Δ "file name"
```

[Function]

- This line is replaced by the specified file name and becomes a processing object as the ST78K0S source program.

[Description]

- The "#" character must always be entered at the start of the symbol, except when starting with a white space or a horizontal tab.
- This directive can be entered in any line in the source program.
- An include directive cannot be entered in an include file. In other words, nesting of include directives is not allowed.
- Input source file names specified at startup, output file names, and error file names cannot be specified as the file name in this directive.
- Drive and directory names can be entered before file names. If no drive or directory is entered, processing assumes that the include file belongs to the current drive and current directory.
- The "-i" option can be used to specify a drive and directory for the include file when the ST78K0S is activated.

[Use examples]

Output source	Input source
MOV A , #08H MOV B , #0AH	#include "sample.inc" A = SYM1 ; #define SYM1 #08H B = SYM2 ; #define SYM2 #0AH

#DEFCALLT

(4) CALLT replacement directive (#defcallt)

[Description format]

```
[ Δ ] # [ Δ ] defcallt Δ CALLT table label
[ Δ ] CALL Δ ! label
[ Δ ] # [ Δ ] endcallt
```

[Function]

- The CALL instruction for a registered label is replaced by a CALLT instruction and is output to a secondary file.

[Description]

- This directive defines labels that can be registered to the CALLT table, as opposed to the CALL instructions that are entered into the source program. All of the CALL instructions for these defined labels are replaced by CALLT labels.
- This directive can be defined up to 32 times. An error message is output during the 33rd definition, and the definition is ignored as processing continues.
- If the same pattern is defined twice, an error message is output and the second definition is ignored as processing continues.

[Use examples]

Output source	Input source
<pre>MOV R0 , #0 CALLT [@ABC] CALL !LABEL CALL_T CSEG AT 40H @ABC : DW ABC</pre>	<pre>#DEFCALLT @ABC CALL !ABC #ENDCALLT R0 = #0 CALL !ABC CALL !LABEL CALL_T CSEG AT 40H @ABC : DW ABC</pre>

CHAPTER 6 CONTROL INSTRUCTIONS

This chapter describes structured assembler control instructions. Control instructions provide detailed instructions for the structured assembler's operations.

6.1 Overview

Control instructions, which are entered into the source program, set various directives that the ST78K0S requires to execute a series of processes.

Entering control instructions saves the time that would otherwise be required for specifying options when activating a program.

6.2 Assembler Control Instructions

First, it must be determined whether or not each assembler control instruction can be entered in a module header.

If there is an assembler control instruction that cannot be entered in a module header, subsequent processing proceeds as the module body. If an assembler control instruction that can only be entered in a module header is instead entered in a module body, an error message is output and processing is aborted.

This preprocessor does not confirm the accuracy of parameter specifications except for processor type specification control instructions (\$PROCESSOR, \$PC), and kanji code specification control instructions (\$KANJI CODE). For description of the description format for other control instructions, see the RA78K0S Assembler Package Language User's Manual.

Table 6-1 lists control instructions that can be entered only in module headers.

Table 6-2 lists control instructions that are recognized as the module body.

Table 6-1 Control Instructions that Can Be Entered Only in Module Headers

Control instruction
[Δ] \$ [Δ] PROCESSOR [Δ] ([Δ] model name [Δ])
[Δ] \$ [Δ] PC ([Δ] model name [Δ])
[Δ] \$ [Δ] DEBUG
[Δ] \$ [Δ] DG
[Δ] \$ [Δ] NODEBAG
[Δ] \$ [Δ] NODG
[Δ] \$ [Δ] DEBUGA
[Δ] \$ [Δ] NODEBAGA
[Δ] \$ [Δ] XREF
[Δ] \$ [Δ] XR
[Δ] \$ [Δ] NOXREF
[Δ] \$ [Δ] NOXR
[Δ] \$ [Δ] TITLE [Δ] ([Δ] 'title string' [Δ])
[Δ] \$ [Δ] TT [Δ] ([Δ] 'title string' [Δ])
[Δ] \$ [Δ] SYMLIST
[Δ] \$ [Δ] NOSYMLIST
[Δ] \$ [Δ] FORMFEED
[Δ] \$ [Δ] NOFORMFEED
[Δ] \$ [Δ] WIDTH [Δ] ([Δ] constant [Δ])
[Δ] \$ [Δ] LENGTH [Δ] ([Δ] constant [Δ])
[Δ] \$ [Δ] TAB [Δ] ([Δ] constant [Δ])
[Δ] \$ [Δ] KANJICODE Δ kanji code

Table 6-2 Control Instructions that Are Recognized as the Module Body

Control instruction
[Δ]\$[Δ] INCLUDE [Δ] ([Δ] file name [Δ])
[Δ]\$[Δ] IC ([Δ] file name [Δ])
[Δ]\$[Δ] EJECT
[Δ]\$[Δ] EJ
[Δ]\$[Δ] LIST
[Δ]\$[Δ] LI
[Δ]\$[Δ] NOLIST
[Δ]\$[Δ] NOLI
[Δ]\$[Δ] GEN
[Δ]\$[Δ] NOGEN
[Δ]\$[Δ] COND
[Δ]\$[Δ] NOCOND
[Δ]\$[Δ] SUBTITLE [Δ] ([Δ] 'character string' [Δ])
[Δ]\$[Δ] ST [Δ] ([Δ] 'character string' [Δ])
[Δ]\$[Δ] SET [Δ] ([Δ] switch name [[Δ]: [Δ] switch name ... [Δ])
[Δ]\$[Δ] RESET [Δ] ([Δ] switch name [[Δ]: [Δ] switch name ... [Δ])
[Δ]\$[Δ] IF [Δ] ([Δ] switch name [[Δ]: [Δ] switch name ... [Δ])
[Δ]\$[Δ] _IF Δ conditional expression
[Δ]\$[Δ] ELSEIF [Δ] ([Δ] switch name [[Δ]: [Δ] switch name ... [Δ])
[Δ]\$[Δ] _ELSEIF Δ conditional expression
[Δ]\$[Δ] ELSE
[Δ]\$[Δ] ENDIF

6.3 Control Instruction Functions

The various functions of control instructions are listed in [Table 6-3](#) below.

Table 6-3 Control Instruction List

Type of control instruction	Control instruction
Processor type specification instruction	\$PROCESSOR / \$PC
Kanji code specification control instructions	\$KANJI CODE

The functions of these three types of control instructions are described below.

\$PROCESSOR / \$PC

(1) Processor type specification instruction (\$PROCESSOR / \$PC)

[Description format]

```
[ Δ ] $ [ Δ ] PROCESSOR [ Δ ] ( [ Δ ] model name [ Δ ] )
[ Δ ] $ [ Δ ] PC [ Δ ] ( [ Δ ] model name [ Δ ] ) ; Abbreviated form
```

[Function]

- This control instruction specifies the model in the source module that is the object for assembly.

[Description]

- Although this control instruction specifies the model that is the object for assembly by the assembler, it can also be used to specify the model that is the object for the structured assembler.
- If the specified model differs from that specified via the "-c" option, the model specified via the "-c" option takes priority. When such a conflict arises, a warning message is output. The "\$" in the input source file's control instruction is replaced by a ";" in the secondary source file that is output, and the model specified via an option is output as the processor model specification control instruction. No message is output if the same model name is specified by the "-c" option. If there is no specification via the "-c" option, the specification must be entered at the start of the source module (not including spaces or comments).
- An error occurs when this control instruction is entered more than once.
- An error occurs if neither this control instruction nor the "-c" option is used to specified a model name.
- An error occurs if this control instruction is entered anywhere other than in the module header.

[Code example]

```
$PROCESSOR ( P9014 )
$PC ( P9014 )
```

\$KANJI CODE

(2) Kanji code specification control instruction (\$KANJI CODE)

[Description format]

```
[ Δ ] $ [ Δ ] KANJI CODE Δ kanji code
```

- Default assumption

Windows : \$KANJI CODE SJIS

Solaris : \$KANJI CODE EUC

[Function]

- The kanji codes used in comments are interpreted as follows.

Table 6-4 Interpretation of Kanji Code

Kanji code	Interpretation
SJIS	Interpreted as SHIFT-JIS code
EUC	Interpreted as EUC code
NONE	Not interpreted as kanji code

[Description]

- This control instruction can be entered in the module header section of an input source file.
- An error occurs if this control instruction is entered anywhere other than in the module header.
- If this control instruction is entered more than once, the most recent one takes priority.
- This preprocessor outputs the specified control instruction to a secondary source file.

SJIS : \$KANJI CODE SJIS

EUC : \$KANJI CODE EUC

NONE : \$KANJI CODE NONE

If the same control instruction is entered in a secondary source file, the control instruction is not output.

However, error checking is performed.

- Kanji code specifications are ranked in terms of priority as follows.
 1. Specification of `-zs/-ze/-zn` option
 2. Specification of the kanji code specification control instruction (\$KANJI CODE)
 3. Specification of the environmental variable LANG78K
 4. Default specification of each OS

[Code example]

```
$KANJI CODE SJIS
```

APPENDIX A SYNTAX LISTS

Table A-1 Control Statements

Control statement	Coding format
if statement if ~ elseif ~ else ~ endif	<pre> if (conditional expression 1) [(register name)] if block elseif (conditional expression 2) [(register name)] elseif block else else block endif </pre>
switch statement switch ~ case ~ default ~ ends	<pre> switch (symbol) [(register name)] case constant 1 : case1 block case constant 2 : case2 block : case constant N : caseN block default : default block ends </pre>
for statement for ~ next	<pre> for (expression ; conditional expression ; expression) [(register name)] Instruction group next </pre>
while statement while ~ endw	<pre> while (conditional expression) [(register name)] Instruction group endw </pre>
until statement repeat ~ until	<pre> repeat Instruction group until (conditional expression) [(register name)] </pre>
break statement break	<pre> break </pre>
continue statement continue	<pre> continue </pre>
goto statement goto	<pre> goto label </pre>
if_bit statement if_bit ~ elseif_bit ~ else ~ endif	<pre> if_bit (conditional expression 1) [(register name)] if_bit block elseif_bit (conditional expression 2) [(register name)]] elseif_bit block else else block endif </pre>
while_bit statement while_bit ~ endw	<pre> while_bit (conditional expression) [(register name)] Instruction group endw </pre>

Table A-1 Control Statements

Control statement	Coding format
until_bit statement repeat ~ until_bit	repeat Instruction group until_bit (conditional expression) [(register name)]

Table A-2 Conditional Expressions

Conditional expression	Coding format	Function
Equal (==)	$\alpha == \beta$	True when $\alpha = \beta$, false when $\alpha \neq \beta$
NotEqual (!=)	$\alpha != \beta$	True when $\alpha \neq \beta$, false when $\alpha = \beta$
LessThan (<)	$\alpha < \beta$	True when $\alpha < \beta$, false when $\alpha \geq \beta$
GreaterThan (>)	$\alpha > \beta$	True when $\alpha > \beta$, false when $\alpha \leq \beta$
GreaterEqual (>=)	$\alpha \geq \beta$	True when $\alpha \geq \beta$, false when $\alpha < \beta$
LessEqual (<=)	$\alpha \leq \beta$	True when $\alpha \leq \beta$, false when $\alpha > \beta$
FOREVER (forever)	forever	Sets endless loop for loop statement
Positive logic (bit) Bit symbol	Bit symbol	True when value of specified bit symbol is 1
Negative logic (bit) !bit symbol	!bit symbol	True when value of specified bit symbol is 0
Logical AND (&&)	Conditional expression 1 && conditional expression 2	True when both conditional expression 1 and conditional expression 2 are true
Logical OR ()	Conditional expression 1 conditional expression 2	True when either conditional expression 1 or conditional expression 2 is true

Table A-3 Expressions

Expression		Coding format	Function
Assign (=)	Assign	$\alpha = \beta$	$\alpha \leftarrow \beta$
	Assign (with register specification)	$\alpha = \beta (\gamma)$	$(\gamma) \leftarrow \beta, \alpha \leftarrow (\gamma)$
	Sequential assign	$\alpha_1 = \dots = \alpha_n = \beta$	$\alpha_1 \leftarrow \beta, \dots, \alpha_n \leftarrow \beta$
	Sequential assign (with register specification)	$\alpha_1 = \dots = \alpha_n = \beta (\gamma)$	$\gamma \leftarrow \beta, \alpha_1 \leftarrow \gamma, \dots, \alpha_n \leftarrow \gamma$
IncrementAssign (+=)	Increment assignment	$\alpha += \beta$	$\alpha \leftarrow \alpha + \beta$
	Increment assignment (with register specification)	$\alpha += \beta (Register)$	$\gamma \leftarrow \alpha, \gamma \leftarrow \gamma + \beta, \alpha \leftarrow \gamma$
	Increment assignment (with register specification)	$\alpha += \beta , CY$	$\alpha \leftarrow \alpha + \beta, CY$
	Increment assignment (with register specification)	$\alpha += \beta , CY (Register)$	$\gamma \leftarrow \alpha, \gamma \leftarrow \gamma + \beta, CY, \alpha \leftarrow \gamma$
DecrementAssign (-=)	Decrement assignment	$\alpha -= \beta$	$\alpha \leftarrow \alpha - \beta$
	Decrement assignment (with register specification)	$\alpha -= \beta (Register)$	$\gamma \leftarrow \alpha, \gamma \leftarrow \gamma - \beta, \alpha \leftarrow \gamma$
	Decrement assignment (with register specification)	$\alpha -= \beta , CY$	$\alpha \leftarrow \alpha - \beta, CY$
	Decrement assignment (with register specification)	$\alpha -= \beta , CY (Register)$	$\gamma \leftarrow \alpha, \gamma \leftarrow \gamma - \beta, CY, \alpha \leftarrow \gamma$
LogicalANDAssign (&=)	Logical AND assignment	$\alpha \&= \beta$	$\alpha \leftarrow \alpha \cap \beta$
	Logical AND assignment (with register specification)	$\alpha \&= \beta (Register)$	$\gamma \leftarrow \alpha, \gamma \leftarrow \gamma \cap \beta, \alpha \leftarrow \gamma$
LogicalORAssign (=)	Logical OR assignment	$\alpha = \beta$	$\alpha \leftarrow \alpha \cup \beta$
	Logical OR assignment (with register specification)	$\alpha = \beta (Register)$	$\gamma \leftarrow \alpha, \gamma \leftarrow \gamma \cup \beta, \alpha \leftarrow \gamma$
LogicalXORAssign (^=)	Logical XOR assignment	$\alpha ^= \beta$	$\alpha \leftarrow \alpha \wedge \beta$
	Logical XOR assignment (with register specification)	$\alpha ^= \beta (Register)$	$\gamma \leftarrow \alpha, \gamma \leftarrow \gamma \wedge \beta, \alpha \leftarrow \gamma$
RightShiftAssign (>>=)	Right shift (rotate) assignment	$\alpha >>= \beta$	$(\alpha \text{ shifted to right of } \beta \text{ bit})$
	Right shift assignment (with register specification)	$\alpha >>= \beta (Register)$	$\gamma \leftarrow \alpha, (\gamma \text{ shifted to right of } \beta \text{ bit}), \alpha \leftarrow \gamma$
LeftShiftAssign (<<=)	Left shift assignment	$\alpha <<= \beta$	$(\alpha \text{ shifted to left of } \beta \text{ bit})$
	Left shift assignment (with register specification)	$\alpha <<= \beta (Register)$	$\gamma \leftarrow \alpha, (\gamma \text{ shifted to left of } \beta \text{ bit}), \alpha \leftarrow \gamma$
Increment (++)	Increment	$\alpha ++$	$\alpha \leftarrow \alpha + 1$
Decrement (--)	Decrement	$\alpha --$	$\alpha \leftarrow \alpha - 1$

Table A-3 Expressions

Expression		Coding format	Function
Exchange (<->)	Exchange	$\alpha \leftarrow \beta$	$\alpha \leftarrow \alpha \leftarrow \beta$
	Exchange (with register specification)	$\alpha \leftarrow \beta (\gamma)$	$\gamma \leftarrow \alpha, \gamma \leftarrow \gamma \leftarrow \beta, \alpha \leftarrow \gamma$
Set bit (=)	Set bit	$\alpha = 1$	$\alpha \leftarrow 1$
	Set bit (with register specification)	$\alpha = 1 (CY)$	$CY \leftarrow 1, \alpha \leftarrow 1$
	Sequential set bit	$\alpha_1 = \dots = \alpha_n = 1$	$\alpha_n \leftarrow 1, \dots, \alpha_1 \leftarrow 1$
	Sequential set bit (with register specification)	$\alpha_1 = \dots = \alpha_n = 1 (CY)$	$CY \leftarrow 1, \alpha_n \leftarrow 1, \dots, \alpha_1 \leftarrow 1$
Clear bit (=)	Clear bit	$\alpha = 0$	$\alpha \leftarrow 0$
	Clear bit (with register specification)	$\alpha = 0 (CY)$	$CY \leftarrow 0, \alpha \leftarrow 0$
	Sequential clear bit	$\alpha_1 = \dots = \alpha_n = 0$	$\alpha_n \leftarrow 0, \dots, \alpha_1 \leftarrow 0$
	Sequential clear bit (with register specification)	$\alpha_1 = \dots = \alpha_n = 0 (CY)$	$CY \leftarrow 0, \alpha_n \leftarrow 0, \dots, \alpha_1 \leftarrow 0$

Table A-4 Directives

Directive	Coding format
#define Symbol definition directive (#define)	#define symbol character string
#ifdef Conditional processing directive (#ifdef / #else / #endif)	#ifdef symbol text 1 #else text 2 #endif
#include Include directive (#include)	#include "file name"
#defcallt CALLT replacement directive (#defcallt)	#defcallt CALLT table label CALL label #endcallt

Table A-5 Control Instructions

Control Instruction	Coding format
Processor type specification instruction (\$PROCESSOR / \$PC)	\$PROCESSOR (model name)
Kanji code specification control instruction (\$KANJI CODE)	\$KANJI CODE Kanji code

APPENDIX B LISTS OF GENERATED INSTRUCTIONS

Table B-1 Generated Instructions for Comparison Expressions

Comparison expression		Generated instruction	Control statement condition
Equal (==)	$\alpha == \beta$	CMP (W) α , β BNZ \$?LFALSE	lower case letters
		CMP (W) α , β BZ \$?LTRUE BR ?LFALSE ?LTRUE :	upper case letters
	$\alpha == \beta (\gamma)$	MOV (W) γ , α CMP (W) γ , β BNZ \$?LFALSE	lower case letters
		MOV (W) γ , α CMP (W) γ , β BZ \$?LTRUE BR LFALSE ?LTRUE :	upper case letters
NotEqual (!=)	$\alpha != \beta$	CMP (W) α , β BZ \$?LFALSE	lower case letters
		CMP (W) α , β BNZ \$?LTRUE BR ?LFALSE ?LTRUE :	upper case letters
	$\alpha != \beta (\gamma)$	MOV (W) γ , α CMP (W) γ , β BZ \$?LFALSE	lower case letters
		MOV (W) γ , α CMP (W) γ , β BNZ \$?LTRUE BR ?LFALSE ?LTRUE :	upper case letters

Table B-1 Generated Instructions for Comparison Expressions

Comparison expression		Generated instruction	Control statement condition
LessThan (<)	$\alpha < \beta$	CMP (W) α , β BNC $\$?LFALSE$	lower case letters
		CMP (W) α , β BC $\$?LTRUE$ BR $?LFALSE$ $?LTRUE :$	upper case letters
	$\alpha < \beta (\gamma)$	MOV (W) γ , α CMP (W) γ , β BNC $\$?LFALSE$	lower case letters
		MOV (W) γ , α CMP (W) γ , β BC $\$?LTRUE$ BR $?LFALSE$ $?LTRUE :$	upper case letters
GreaterThan (>)	$\alpha > \beta$	CMP (W) α , β BZ $\$?LFALSE$ BC $\$?LFALSE$	lower case letters
		CMP (W) α , β BZ $\$\$ + 4$ BNC $\$?LTRUE$ BR $?LFALSE$ $?LTRUE :$	upper case letters
	$\alpha > \beta (\gamma)$	MOV (W) specified register , α CMP (W) specified register , β BZ $\$?LFALSE$ BC $\$?LFALSE$	lower case letters
		MOV (W) specified register , α CMP (W) specified register , β BZ $\$\$ + 4$ BNC $\$?LTRUE$ BR $?LFALSE$ $?LTRUE :$	upper case letters
GreaterEqual (>=)	$\alpha >= \beta$	CMP (W) α , β BC $\$?LFALSE$	lower case letters
		CMP (W) α , β BNC $\$?LTRUE$ BR $?LFALSE$ $?LTRUE :$	upper case letters
	$\alpha >= \beta (\gamma)$	MOV (W) γ , α CMP (W) γ , β BC $\$?LFALSE$	lower case letters
		MOV (W) γ , α CMP (W) γ , β BNC $\$?LTRUE$ BR $?LFALSE$ $?LTRUE :$	upper case letters

Table B-1 Generated Instructions for Comparison Expressions

Comparison expression		Generated instruction	Control statement condition
LessEqual (<=)	$\alpha \leq \beta$	CMP (W) α , β BZ $$$ + 4$ BNC $$$?LFALSE$	lower case letters
		CMP (W) α , β BZ $$$?LTRUE$ BC $$$?LTRUE$ BR $$$?LFALSE$?LTRUE :	upper case letters
	$\alpha \leq \beta (\gamma)$	MOV (W) specified register , α CMP (W) specified register , β BZ $$$ + 4$ BNC $$$?LFALSE$	lower case letters
		MOV (W) specified register , α CMP (W) specified register , β BZ $$$?LTRUE$ BC $$$?LTRUE$ BR $$$?LFALSE$?LTRUE :	upper case letters

γ : specified register

Table B-2 Generated Instructions for Test Bit Expressions

Test bit expression	Generated instruction	Control statement condition
Bit symbol if_bit (bit symbol) elseif_bit (bit symbol) while_bit (bit symbol) until_bit (bit symbol)	BNC \$?LFALSE	lower case letters (CY)
	BNZ \$?LFALSE	lower case letters (Z)
	BF bit symbol , \$?LFALSE	lower case letters
	BC \$?LTRUE BR ?LFALSE ?LTRUE :	upper case letters (CY)
	BZ \$?LTRUE BR ?LFALSE ?LTRUE :	upper case letters (Z)
	BT bit symbol , \$?LTRUE BR ?LFALSE ?LTRUE :	upper case letters
!bit symbol if_bit (!bit symbol) elseif_bit (!bit symbol) while_bit (!bit symbol) until_bit (!bit symbol)	BC \$?LFALSE	lower case letters (CY)
	BZ \$?LFALSE	lower case letters (Z)
	BT bit symbol , \$?LFALSE	lower case letters
	BNC \$?LTRUE BR ?LFALSE ?LTRUE :	upper case letters (CY)
	BNZ \$?LTRUE BR ?LFALSE ?LTRUE :	upper case letters (Z)
	BF bit symbol , \$?LTRUE BR ?LFALSE ?LTRUE :	upper case letters

Table B-3 Generated Instructions for Logic Expressions

Logic expression		Generated instruction		Control statement condition
Logical AND (&&)	$\alpha == \beta$ &&	CMP (W)	α , β	lower case letters
		BNZ	$\$?LFALSE$	
	$\alpha == \beta$ &&	CMP (W)	α , β	upper case letters
		BZ	$\$?LTRUE$	
		BR	$?LFALSE$	
	$\alpha != \beta$ &&	?LTRUE :		
		CMP (W)	α , β	lower case letters
		BZ	$\$?LFALSE$	
	$\alpha != \beta$ &&	CMP (W)	α , β	upper case letters
		BNZ	$\$?LTRUE$	
		BR	$?LFALSE$	
	$\alpha != \beta$ &&	?LTRUE :		
		CMP (W)	α , β	lower case letters
		BNC	$\$?LFALSE$	
$\alpha < \beta$ &&	CMP (W)	α , β	upper case letters	
	BC	$\$?LTRUE$		
	BR	$?LFALSE$		
$\alpha < \beta$ &&	?LTRUE :			
	CMP (W)	α , β	lower case letters	
	BZ	$\$?LFALSE$		
$\alpha > \beta$ &&	BC	$\$?LFALSE$		
	CMP (W)	α , β	upper case letters	
	BZ	$\$\$ + 4$		
BNC	$\$?LTRUE$			
$\alpha > \beta$ &&	BR	$?LFALSE$		
	?LTRUE :			
	CMP (W)	α , β	lower case letters	
$\alpha >= \beta$ &&	BC	$\$?LFALSE$		
	CMP (W)	α , β	upper case letters	
	BNC	$\$?LTRUE$		
BR	$?LFALSE$			
$\alpha >= \beta$ &&	?LTRUE :			
	CMP (W)	α , β	lower case letters	
	BZ	$\$\$ + 4$		
$\alpha <= \beta$ &&	BNC	$\$?LFALSE$		
	CMP (W)	α , β	upper case letters	
	BZ	$\$?LTRUE$		
BC	$\$?LTRUE$			
$\alpha <= \beta$ &&	BR	$?LFALSE$		
	?LTRUE :			

Table B-3 Generated Instructions for Logic Expressions

Logic expression		Generated instruction	Control statement condition
Logical AND (&&)	CY &&	BNC \$?LFALSE	lower case letters
		BC \$?LTRUE BR ?LFALSE ?LTRUE :	upper case letters
	Z &&	BNZ \$?LFALSE	lower case letters
		BZ \$?LTRUE BR ?LFALSE ?LTRUE :	upper case letters
	bit symbol &&	BF bit symbol , \$?LFALSE	lower case letters
		BT bit symbol , \$?LTRUE BR ?LFALSE ?LTRUE :	upper case letters
	!CY &&	BC \$?LFALSE	lower case letters
		BNC \$?LTRUE BR ?LFALSE ?LTRUE :	upper case letters
	!Z &&	BZ \$?LFALSE	lower case letters
		BNZ \$?LTRUE BR ?LFALSE ?LTRUE :	upper case letters
	!bit symbol &&	BT bit symbol , \$?LFALSE	lower case letters
		BF bit symbol , \$?LTRUE BR ?LFALSE ?LTRUE :	upper case letters

Table B-3 Generated Instructions for Logic Expressions

Logic expression		Generated instruction		Control statement condition
Logical OR ()	$\alpha == \beta$	CMP (W) BZ	α , β \$?LFALSE	none
	$\alpha != \beta$	CMP (W) BNZ	α , β \$?LFALSE	
	$\alpha < \beta$	CMP (W) BC	α , β \$?LFALSE	
	$\alpha > \beta$	CMP (W) BZ BNC	α , β \$?LFALSE \$?LFALSE	
	$\alpha >= \beta$	CMP (W) BNC	α , β \$?LFALSE	
	$\alpha <= \beta$	CMP (W) BZ BC	α , β \$?LFALSE \$?LFALSE	
	CY	BC	\$?LFALSE	
	Z	BZ	\$?LFALSE	
	bit symbol	BT	bit symbol , \$?LFALSE	
	!CY	BNC	\$?LFALSE	
	!Z	BNZ	\$?LFALSE	
	!bit symbol	BF	bit symbol , \$?LFALSE	

Table B-4 Expressions

Expression		Generated instruction	
Assign (=)	$\alpha = \beta$	MOV α_1, β	
		MOVW α_1, β	
		BNC ?L1 SET1 α BR ?L2	
		?L1 : CLR1 α ?L2 :	
	$\alpha = \beta(\gamma)$	MOV γ, β MOV α_1, γ	
		MOVW γ, β MOVW α_1, γ	
		BF $\beta, ?L1$ SET1 α BR ?L2	
		?L1 : CLR1 α ?L2 :	
	IncrementAssign (+=)	$\alpha += \beta$	ADD α, β
			ADDW α, β
$\alpha += \beta(\gamma)$		MOV γ, α ADD γ, β MOV α, γ	
		MOVW γ, α ADDW γ, β MOVW α, γ	
$\alpha += \beta, CY$		ADDC α, β	
$\alpha += \beta, CY(\gamma)$	MOV γ, α ADDC γ, β MOV α, γ		
DecrementAssign (-=)	$\alpha -= \beta$	SUB α, β	
		SUBW α, β	
	$\alpha -= \beta(\gamma)$	MOV γ, α SUB γ, β MOV α, γ	
		MOVW γ, α SUBW γ, β MOVW α, γ	
	$\alpha -= \beta, CY$	SUBC α, β	
$\alpha -= \beta, CY(\gamma)$	MOV γ, α SUBC γ, β MOV α, γ		

Table B-4 Expressions

Expression		Generated instruction
LogicalANDAssign (&=)	$\alpha \&= \beta$	AND α , β
		BNC ?L1 BF $\beta , ?L1$ SET1 CY BR ?L2 ?L1 : CLR1 CY ?L2 :
	$\alpha \&= \beta (\gamma)$	MOV γ , α AND γ , β MOV α , γ
		BF $\alpha , ?L1$ BF $\beta , ?L1$ SET1 α BR ?L2 ?L1 : CLR1 α ?L2 :
LogicalORAssign (=)	$\alpha = \beta$	OR α , β
		BC ?L1 BF $\beta , ?L2$?L1 : SET1 CY BR ?L3 ?L2 : CLR1 CY ?L3 :
		MOV γ , α OR γ , β MOV α , γ
	$\alpha = \beta (\gamma)$	BT $\alpha , ?L1$ BF $\beta , ?L2$?L1 : SET1 α BR ?L3 ?L2 : CLR1 α ?L3 :

Table B-4 Expressions

Expression		Generated instruction
LogicalXORAssign (^=)	$\alpha \wedge = \beta$	XOR α, β BNC ?L1 BF $\beta, ?L2$?L1 : BC ?L3 BF $\beta, ?L3$?L2 : SET1 CY BR ?L4 ?L3 : CLR1 CY ?L4 :
	$\alpha \wedge = \beta (\gamma)$	MOV γ, α XOR γ, β MOV α, γ BF $\alpha, ?L1$ BF $\beta, ?L2$?L1 : BT $\alpha, ?L3$ BF $\beta, ?L3$?L2 : SET1 α BR ?L4 ?L3 : CLR1 α ?L4 :
RightShiftAssign (>>=)	$\alpha >> = \beta$	ROR A, 1 : AND A, #0FFH SHR β
	$\alpha >> = \beta (\gamma)$	MOV A, α ROR A, 1 : AND A, #0FFH SHR β MOV α, A
LeftShiftAssign (<<=)	$\alpha << = \beta$	ROL A, 1 : AND A, #LOW (0FFH SHL β)
	$\alpha << = \beta (\gamma)$	MOV A, α ROL A, 1 : AND A, #LOW (0FFH SHL β) MOV α, A
Increment (++)	$\alpha ++$	INC α INCW α
Decrement (--)	$\alpha --$	DEC α DECW α

Table B-4 Expressions

Expression		Generated instruction
Exchange (<->)	$\alpha \leftrightarrow \beta$	XCH α, β
		XCHW α, β
	$\alpha \leftrightarrow \beta (\gamma)$	MOV γ, α XCH γ, β MOV α, γ
		MOVW γ, α XCHW γ, β MOVW α, γ
Set bit (=)	$\alpha = 1$	SET1 α_1
	$\alpha = 1 (CY)$	SET1 CY SET1 α_1
Clear bit (=)	$\alpha = 0$	CLR1 α_1
	$\alpha = 0 (CY)$	CLR1 CY CLR1 α_1

APPENDIX C INDEX

Symbols

#DEFCALLT ... 143
#DEFINE ... 138
#ELSE ... 140
#ENDIF ... 140
#IFDEF ... 31, 140
#INCLUDE ... 31, 142
\$KANJI CODE ... 149
\$PC ... 148
\$PROCESSOR ... 148

A

Assembler control instruction ... 24
Assembler operator ... 23
Assembly language ... 18
Assign ... 99
Assignment statement ... 96, 99

B

Bit manipulation statement ... 96, 98, 130
break ... 57
Byte access ... 22
Byte symbol ... 17

C

CALLT replacement directive ... 143
Character set ... 20
Clear bit ... 133
Comment ... 31
Comment statement ... 30
Comparison expression ... 61
Conditional branch ... 37
Conditional expression ... 18
Conditional loop ... 47
Conditional processing directive ... 140
Constant ... 22
continue ... 58
Control statement ... 18, 23, 31
Count statement ... 96, 97, 123

D

Decrement ... 125
DecrementAssign ... 107
Directive ... 23, 31, 136

E

Equal ... 63
Error message ... 30
Exchange ... 127
Exchange statement ... 96, 98, 127
Expression ... 22
Expression statement ... 18, 31

F

for ... 47
FOREVER (forever) ... 81

G

goto ... 59
GreaterEqual ... 75
GreaterThan ... 72

I

Identifier ... 22
if ... 37
if_bit ... 40
Illegal character ... 22, 29
Include directive ... 142
Increment ... 123
IncrementAssign ... 104
Invalid character ... 22

K

Kanji code specification control instruction ... 149

L

Label ... 17, 25
LeftShiftAssign ... 121
LessEqual ... 78
LessThan ... 69
Letter ... 20
Line feed ... 30
Logical AND ... 91
Logical operation ... 90
Logical OR ... 94
LogicalANDAssign ... 110
LogicalORAssign ... 113
LogicalXORAssign ... 116
Lower case letter ... 20, 32

M

Module body ... 145
Module header ... 30, 145

N

Negative logic (bit) ... 87
NotEqual ... 66
Numeral ... 20

O

Operator ... 23

P

Positive logic (bit) ... 84
Processor type specification instruction ... 148

R

Register ... 24
RightShiftAssign ... 119

S

Set bit ... 130
Source program ... 136, 144
Special character ... 20
ST78K0S ... 13
Structured assembler preprocessor ... 13
Structured assembly language ... 18
switch ... 43
Symbol ... 22
Symbol definition directive ... 138

T

Test bit expression ... 83

U

until ... 54
until_bit ... 56
Upper case letter ... 20, 32
User symbol ... 22

W

while ... 50
while_bit ... 52
Word access ... 22
Word symbol ... 17

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